

ANTICHAINS IN THE HOMOMORPHISM ORDER OF GRAPHS

D. DUFFUS, P. L. ERDŐS, J. NEŠETŘIL, AND L. SOUKUP

ABSTRACT. Let \mathbb{G} and \mathbb{D} , respectively, denote the partially ordered sets of homomorphism classes of finite undirected and directed graphs, respectively, both ordered by the homomorphism relation. Order theoretic properties of both have been studied extensively, and have interesting connections to familiar graph properties and parameters. In particular, the notion of a duality is closely related to the idea of splitting a maximal antichain. We construct both splitting and non-splitting infinite maximal antichains in \mathbb{G} and in \mathbb{D} . The splitting maximal antichains give infinite versions of dualities for graphs and for directed graphs.

1. INTRODUCTION

For any fixed type of finite relational structure, homomorphisms induce an ordering of the set of all structures. In particular, given two graphs [respectively, directed graphs] G and H write $G \leq H$ or $G \rightarrow H$ provided that there is a homomorphism from G to H , that is, a map $f : V(G) \rightarrow V(H)$ such that for all $\{x, y\} \in E(G)$, $\{f(x), f(y)\} \in E(H)$ [respectively, for all $\langle x, y \rangle \in E(G)$, $\langle f(x), f(y) \rangle \in E(H)$]. Then the relation \leq is a quasi-order and so it induces an equivalence relation: we say that G and H are *homomorphism-equivalent* or *hom-equivalent* and write $G \sim H$ if and only if $G \leq H$ and $H \leq G$. The *homomorphism posets* \mathbb{G} and \mathbb{D} are the partially ordered sets of all equivalence classes of finite undirected and directed graphs, respectively, ordered by the \leq . We will often abuse notation by replacing the classes that comprise \mathbb{G} and \mathbb{D} with their members.

These partially ordered sets are of significant intrinsic interest and are useful tools in the study of graph and directed graph properties. For instance, it is easily seen that both are countable distributive lattices: the supremum, or join, of any pair is their disjoint sum, and the infimum, or meet, is their categorical or relational product. Both \mathbb{G} and \mathbb{D} are “predominantly” dense – the former shown by Welzl [20] and the latter, by Nešetřil and Tardif [16]. Both also embed all countable partially ordered sets – see [19] for a presentation. Basic order-theoretic

Date: February 25, 2008.

The second author was partly supported by Hungarian NSF, under contract Nos. T 37758, T37846, NK62321, AT048826. The fourth author was supported by Hungarian NSF, under contract No. 61600.

THIS LONG VERSION CONTAINS THE PROOFS WHICH ARE MISSING FROM THE SUBMITTED VERSION

properties, such as the existence of suprema and infima for several natural families in \mathbb{G} , are considered in [13].

The maximal chains and antichains of an ordered set are subobjects of interest. In this case, maximal antichains are particularly relevant because of their relationship to the notion of a *homomorphism duality*, introduced by Nešetřil and Pultr [14]: say that an ordered pair $\langle F, D \rangle$ of graphs, or directed graphs, is a *duality pair* if

$$(1) \quad F \rightarrow = \nrightarrow D$$

where $F \rightarrow = \{G : F \rightarrow G\}$ and $\nrightarrow D = \{G : G \nrightarrow D\}$. Equivalently, the set of all structures is partitioned by the *upset* [or *final segment*] $F \rightarrow$ and the *downset* [or *initial segment*] $\rightarrow D$. [Here we also use the other common notation F^\uparrow and D^\downarrow for upsets and downsets, respectively.]

One important motivation for consideration of duality pairs is that of an “obstruction” to a graph property. For instance, the possibility of a homomorphism of a graph G to K_2 , a 2-coloring, is obstructed by the existence of a homomorphism of some odd cycle to G . While there are no nontrivial duality pairs in \mathbb{G} , in \mathbb{D} , each tree can play the role of F in (1). In fact, in [16], Nešetřil and Tardif obtain a correspondence between duality pairs and gaps in the homomorphism order for general relational structures. They use this to characterize duality pairs and generalize this by describing exactly when the left handside of (1) can be replaced by a finite union of final segments. They further note in [17] that the 2-element maximal antichains in \mathbb{D} are exactly the duality pairs $\langle F, D \rangle$ where F is a tree and D is its dual.

Foniok, Nešetřil and Tardif [10] are concerned with the most general circumstance. Let \mathcal{F} and \mathcal{D} both be finite antichains of structures of fixed type Δ . Call $\langle \mathcal{F}, \mathcal{D} \rangle$ a *generalized duality* if

$$(2) \quad \bigcup_{F \in \mathcal{F}} F \rightarrow = \bigcap_{D \in \mathcal{D}} \nrightarrow D$$

Equivalently, with \mathbb{S} denoting the homomorphism poset of Δ -structures, \mathbb{S} is partitioned by

$$(3) \quad \mathbb{S} = \left(\bigcup_{F \in \mathcal{F}} F \rightarrow \right) \cup \left(\bigcup_{D \in \mathcal{D}} \rightarrow D \right).$$

The generalized dualities are characterized in [10]. They also show that when Δ consists of one k -ary relation, which contains the graph cases, every finite maximal antichain in the lattice of Δ -structures is of the form $\mathcal{F} \cup \mathcal{D}$. Conversely, for all but three exceptional cases, the generalized dualities $\langle \mathcal{F}, \mathcal{D} \rangle$ yield a maximal antichain $\mathcal{F} \cup \mathcal{D}$.

It is quite natural to ask, in more general circumstances, if maximal antichains possess these sorts of partitions. Indeed, Ahlswede, Erdős and Graham [1] introduced the notion of “splitting” a maximal antichain. Say that a maximal antichain A of a poset P *splits* if A can be partitioned into two subsets B and C such that $P = B^\uparrow \cup C^\downarrow$; say that P has the *splitting property* if all of its maximal antichains

split. They obtained sufficient conditions for the splitting property, from which they proved, in particular, that all finite Boolean lattices possess it. The property is also a useful tool in combinatorial investigations of posets, particularly distributive lattices; see, for instance [2, 3]. It is also a natural notion for infinite posets; see [4, 8, 9].

The correspondence between generalized dualities and maximal antichains obtained in [10] and the partition in (3) demonstrate that for $\Delta = (k)$, essentially all finite maximal antichains in the lattice \mathbb{S} of Δ -structures split.

This paper is motivated by two goals. First, we would like to obtain general order theoretic conditions on countable posets that ensure antichains split and, thereby, obtain some of the duality results that had been restricted to finite maximal antichains, as described above. See Section 4 for applications to \mathbb{G} and Section 5 for results on \mathbb{D} . Second, we obtain splitting and non-splitting results for infinite maximal antichains; in particular, these results underscore differences between the structures \mathbb{G} and \mathbb{D} . The necessary results on splitting and related notions are given in Section 3, which is preceded in Section 2 by a directed version of what is known as the Sparse Incomparability Lemma.

In addition to the selected papers cited in this section, we refer the reader to the book [11] by Hell and Nešetřil that is devoted to graph homomorphisms. Chapter 3 gives a thorough introduction and many of the key results on maximal antichains and dualities in \mathbb{G} and \mathbb{D} .

2. A DIRECTED SPARSE INCOMPARABILITY LEMMA

Recall that the *girth* of a graph, $\text{girth}(G)$, is the length of a shortest cycle contained in the graph. In case G is directed, its girth is that of the underlying undirected graph, that is, of the symmetric version of G . In one of the first applications of the probabilistic method, in 1959 Paul Erdős [5] showed the existence of graphs with independently prescribed girth and chromatic number. More precisely, for all natural numbers k and ℓ there is a graph G such that $\chi(G) > k$ and $\text{girth}(G) > \ell$.

Based on another probabilistic argument due to Erdős and Hajnal [6], Nešetřil and Rödl [15] obtained an interesting generalization, referred to as the “Sparse Incomparability Lemma”: for every pair of graphs H and G such that $G \rightarrow H$ but $H \not\rightarrow G$, and for every positive integer ℓ there exists a graph H' with $\text{girth}(H') > \ell$ such that $H' \rightarrow H$ and $H' \not\rightarrow G$.

Here is a formulation from which the Sparse Incomparability Lemma follows, itself a special case of a more far-reaching generalization.

Theorem 2.1 (Nešetřil-Zhu [18]). *For every graph H and for every positive integers k and ℓ there exists a graph G with the following properties:*

- (i): $\text{girth}(G) > \ell$, and
- (ii): *for every graph H_0 with at most k vertices, $G \rightarrow H_0$ if and only if $H \rightarrow H_0$.*

Here, we require a directed graph version of Theorem 2.1. The following is a special case of a Sparse Incomparability Lemma for finite relational structures [12].

Theorem 2.2 (Directed Sparse Incomparability Lemma). *For each directed graph $H = (W, F)$ and for all integers $m, \ell \in \mathbb{N}$ there is a directed graph H' such that*

- (i): $\text{girth}(H') > \ell$,
- (ii): *for each directed graph G with $|V(G)| < m$ we have $H' \rightarrow G$ if and only if $H \rightarrow G$, and*
- (iii): *H and H' have the same numbers of connected components. In particular, if H is connected then so is H' .*

Regarding the proof of Theorem 2.2, there are both probabilistic and deterministic arguments available. For instance, it is straightforward to adapt the probabilistic proof of Nešetřil-Rödl. We found an alternative approach based on what appears to be a new graph parameter. Here is a brief outline of the argument.

Given a graph $G = (V, E)$ let the *bipartite stability number* $\alpha_b(G)$ be the maximum integer β such that:

$$\exists A, B \in [V]^\beta \text{ with } A \cap B = \emptyset \text{ and no edge between } A \text{ and } B.$$

Clearly $\alpha_b(G) \geq \alpha(G)/2$, where $\alpha(G)$ denotes the usual stability or independence number of G . The following result is obtained by adjusting the Erdős-Rényi proof [7] that there are graphs of large girth and small independence number.

Lemma 2.3. *For all $k, \ell \in \mathbb{N}$ and for all but finitely many $n \in \mathbb{N}$ there exists a connected graph $G' = (V, E)$ with $|V| = n$, $\text{girth}(G') > \ell$ and $\alpha_b(G') < n/k$.*

Let H , m and ℓ be as in the statement of Theorem 2.2. Let $k = 3m|W|$ and $n = kj$ for sufficiently large j . By Lemma 2.3, there exists a graph $G' = (V, E)$ such that

- $V = W \times [3mj]$
- $\text{girth}(G') > \ell$
- $\alpha_b(G') < n/k = j$

In effect, we “blow up” each vertex of H into a class of $3mj$ vertices.

Define a directed graph $H^* = (V, E^*)$ as follows: if $\langle h, i \rangle, \langle h', i' \rangle \in V$ then $\langle \langle h, i \rangle, \langle h', i' \rangle \rangle \in E^*$ if and only if $(\langle h, i \rangle, \langle h', i' \rangle) \in E$ and $\langle h, h' \rangle \in F$.

One now argues that if H is connected then H^* has a large enough connected component that satisfies (1), (2), and (3) of the theorem.

ADDENDUM

Here we give one possible proof of Theorem 2.2, for the sake of completeness. One can easily come up with shorter probabilistic proof.

Theorem A (Directed Sparse Incomparability Lemma) *For each directed graph $H = (W, F)$ and for all integers $m, \ell \in \mathbb{N}$ there is a directed graph H' such that*

- (1) $\text{girth}(H') > \ell$
- (2) *for each directed graph G with $|V(G)| < m$ we have $H' \rightarrow G$ if and only if $H \rightarrow G$.*
- (3) $\chi(H) = \chi(H')$.
- (4) *H and H' have the same numbers of connected components. Especially, if H is connected then so is H' .*

The proof is based on a Paul Erdős-type result (see Theorem 2.3). At first we need one more definition:

Definition B Given a graph $G = (V, E)$ let the *bipartite stability number*

$$\alpha_b(G) = \max \left\{ \beta \in \mathbb{N} : \exists A, B \in [V]^\beta \ A \cap B = \emptyset \right. \\ \left. \text{and there is no edge between } A \text{ and } B \right\}.$$

Clearly $\alpha_b(G) \geq \alpha(G)/2$. We need:

Theorem C *For all $k, \ell \in \mathbb{N}$ and for all but finitely many $n \in \mathbb{N}$ there exists a connected graph $G' = (V, E)$ with $|V| = n$, $\text{girth}(G') > \ell$ and $\alpha_b(G') < n/k$.*

Proof. Fix $\theta < 1/\ell$ and let $G \sim G(n, p)$ with $p = n^{\theta-1}$. Let X be the number of circuits of size at most ℓ . Then:

$$E[X] = \sum_{i=3}^{\ell} \frac{\binom{n}{i}}{2i} p^i \leq \sum_{i=3}^{\ell} \frac{n^{\theta i}}{2i} = o(n).$$

as $\theta\ell < 1$. In particular,

$$(4) \quad \Pr[X \geq n/2k] = o(1).$$

Set $x = (3/p) \ln n$ so that

$$(5) \quad \Pr[\alpha_b(G) \geq x] \leq \binom{n}{x} \binom{n-x}{x} (1-p)^{x^2} < n^{2x} [e^{-px}]^x = [n^2 e^{-px}]^x = \\ [e^{2 \ln n - p(3/p) \ln n}]^x = [e^{-\ln n}]^x = e^{-\ln n (3/p) \ln n} = e^{-3n^{1-\theta} \ln^2 n} = o(1).$$

Since $p = n^{\theta-1} > \frac{\log n}{n}$ therefore by the well known Erdős-Rényi theorem ([7]) the graph almost surely connected, therefore

$$(6) \quad \Pr[G \text{ is not connected}] = o(1).$$

Let n be sufficiently large so that

- (1) the probabilities in (4), (5) and (6) are all less than $1/3$,
- (2) $x = \lceil (3/p) \ln n \rceil = \lceil 3n^{1-\theta} \ln n \rceil < n/2k$.

Then there is a specific G with less than $n/2k$ cycles of length $\leq \ell$ and with $\alpha_b(G) < n/2k$. Remove, one by one, one edge from each cycle of G of length at most ℓ . This new graph G' has a girth greater than ℓ . Furthermore if $\alpha_b(G) = \alpha$ and we delete one edge from it then $\alpha(G \setminus e) \leq \alpha + 1$. Therefore $\alpha_b(G') \leq n/2k + n/2k \leq n/k$, as required. Finally the described operation clearly preserve the connectedness of G . \square

Denote by \vec{K}_n the full directed graph on n points: $\vec{K}_n = (\{0, \dots, n-1\}, E)$, where $E = \{\langle i, j \rangle : i \neq j < n\}$. Clearly $\chi(H) \leq n$ if and only if $H \rightarrow \vec{K}_n$.

Proof. of Theorem A: We can assume that $m > |W|$.

Assume first that H is connected.

Let $k = 3m|W|$. Let $n = k \cdot j$ for some large enough j . By Theorem 2.3 we can obtain an undirected $G' = (V, E)$ on n vertices with $\text{girth}(G') > \ell$ and $\alpha_b(G') < n/k = j$. We can assume that $V = W \times 3mj$. For $w \in W$ write $V_w = \{w\} \times mj$. (We will "blow up" each vertex of H into a class of $3mj$ vertices.)

Define a digraph graph $H^* = (V, E^*)$ as follows: if $\langle h, i \rangle, \langle h', i' \rangle \in V$ then $\langle \langle h, i \rangle, \langle h', i' \rangle \rangle \in E^*$ if and only if $(\langle h, i \rangle, \langle h', i' \rangle) \in E$ and $\langle h, h' \rangle \in F$.

One could prove that H^* satisfies (1)–(3), but unfortunately H^* is not necessarily connected. We are going to show that the greatest connected component of H^* works as H' , i.e. it satisfies all the requirements (1)–(4).

Claim D H^* has a connected component of size $\geq n/3$.

Proof. If the Claim fails then there is a partition (X_0, X_1) of V such that $n/3 \leq |X_0|, |X_1| \leq 2n/3$ and there is no edge in H^* between X_0 and X_1 .

For $i < 2$ let $Y_i = \{w \in W : |X_i \cap V_w| \geq n/k\}$. Since $|V_w| = n/|W| \geq 2(n/k)$ and $V_w = (V_w \cap X_0) \cup (V_w \cap X_1)$ we have $Y_0 \cup Y_1 = W$.

For $i < 2$ we have $|X_i| \geq n/3$ so there is $w_i \in W$ such that $X_i \cap V_{w_i} \geq n/(3|W|) \geq n/k$, and so $w_i \in Y_i$.

Since $Y_0 \neq \emptyset \neq Y_1$ and $Y_0 \cup Y_1 = W$ there are non-empty disjoint sets $Z_0 \subset Y_0$ and $Z_1 \subset Y_1$ such that $Z_0 \cup Z_1 = W$.

Since $H = (W, F)$ is connected there is $\langle w, w' \rangle \in F$ such that $w \in Z_i$ and $w' \in Z_{1-i}$. Then $\alpha_b(G') < n/k \leq |X_i \cap V_w|, |X_{1-i} \cap V_{w'}|$ holds which implies that there is $a \in X_i \cap V_w$ and $a' \in X_{1-i} \cap V_{w'}$ with $(a, a') \in E$. Then $(a, a') \in E^*$ is an edge between X_i and X_{1-i} . Contradiction. \square

Claim E H^* has an even bigger connected component X_0 , with size

$$|V(X_0)| \geq \left(1 - \frac{1}{3m}\right)n.$$

Proof. Let X_0 be the greatest connected component and $X_1 = V \setminus X_0$.

For $i < 2$ let $Y_i = \{w \in W : |X_i \cap V_w| \geq n/k\}$. Since $|V_w| = n/|W| \geq 2(n/k)$ and $V_w = (V_w \cap X_0) \cup (V_w \cap X_1)$ we have $Y_0 \cup Y_1 = W$.

If the Claim fails then $|X_1| > n/(3m)$. Thus there is $w_1 \in W$ such that $|X_1 \cap V_{w_1}| \geq (n/(3m))/|W| = n/k$. Thus $Y_1 \neq \emptyset$.

Since $|X_0| \geq n/3$ there is $w_0 \in W$ such that $X_0 \cap V_{w_0} \geq (n/3)/|W| \geq n/k$. Thus $Y_0 \neq \emptyset$ as well.

Since furthermore $Y_0 \cup Y_1 = W$ there are non-empty disjoint sets $Z_0 \subset Y_0$ and $Z_1 \subset Y_1$ such that $Z_0 \cup Z_1 = W$.

Since $H = (W, F)$ is connected there is $\langle w, w' \rangle \in F$ such that $w \in Z_i$ and $w' \in Z_{1-i}$. Then $\alpha_b(G') < n/k \leq |X_i \cap V_w|, |X_{1-i} \cap V_{w'}|$ implies that there is $a \in X_i \cap V_w$ and $a' \in X_{1-i} \cap V_{w'}$ with $(a, a') \in E$. Then $(a, a') \in E^*$ is an edge between X_i and X_{1-i} . Contradiction. \square

Let H' be the greatest connected component of H^* . We show that H' satisfies (1)–(4).

Clearly $\text{girth}(H') \geq \text{girth}(H^*) \geq \text{girth}(G') > \ell$ holds. (Recall if there is no cycle in the graph, then the girth is ∞ .) Hence (1) holds.

Write $V' = V(H')$. To check (2) assume that $f : V' \rightarrow V(G)$ witnesses $H' \rightarrow G$.

Let $w \in W$ be fixed. Then

$$|V_w \cap V'| \geq \left(1 - \frac{1}{3m}\right)n - (|W| - 1)\frac{n}{|W|} = \left(n - \frac{n}{3m}\right) - \left(n - \frac{n}{|W|}\right),$$

therefore $|V_w \cap V'| \geq n/|W| - n/(3m) \geq n/|W| - n/(3|W|) > n/(3|W|)$ (recall that $m \geq |W|$). Furthermore $|\text{range}(f)| \leq V(G) < m$ so there is $v_h \in V(G)$ such that $|V_h \cap V' \cap f^{-1}\{v_h\}| \geq (n/3|W|)/m = n/k$. Let $f'(h) = v_h$. Then f' witnesses $H \rightarrow G$.

To prove (3) it is enough to observe that $H \not\rightarrow \vec{K}_{\chi(H)-1}$, so $H' \not\rightarrow \vec{K}_{\chi(H)-1}$ by (2). Thus $\chi(H') \geq \chi(H)$.

So far we proved the theorem for connected digraphs. Now let H be an arbitrary digraphs and let (H_1, \dots, H_j) be an enumeration of the connected components of H . Applying the theorem for each connected components we get H'_i for each $i = 1, \dots, j$. Then $H' = H_1 + H_2 + \dots + H_j$ works for H . Indeed $\text{girth}(H') = \min_{1 \leq i \leq j} \{\text{girth } H'_i\} > \ell$. (2) holds because $H \rightarrow G$ if and only if we have $H_i \rightarrow G$ for each $1 \leq i \leq j$. Moreover $\chi(H) = \max_{1 \leq i \leq j} \{\chi(H_i)\} = \max_{1 \leq i \leq j} \{\chi(H'_i)\} = \chi(H')$. \square

This approach in Theorem 2.2 gives a new proof to Theorem 2.1 which gives slightly stronger result as well:

Theorem F *For every undirected graph $H = (V, E)$ and for every positive integers k and ℓ there exists a graph G with the following properties:*

- (i) $\text{girth}(G) > \ell$,
- (ii) *For every graph H_0 with at most k vertices, $G \rightarrow H_0$ if and only if $H \rightarrow H_0$,*
- (iii) $\chi(H) = \chi(G)$,
- (iv) *H and G have the same numbers of connected components. Especially, if H is connected then so is G .*

Proof. Let the directed graph $\vec{H} = (V, F)$ defined with $F = \{\langle x, y \rangle, \langle y, x \rangle : (x, y) \in E\}$. The application of Theorem 2.2 provides us the directed graph \vec{G} . Let \vec{G}' be the directed graph derived from \vec{G} with oriented all edges on both directions. It is easy to see that the graph \vec{G}' also satisfies each condition of Theorem 2.2. Indeed, properties (i), (iii) and (iv) are clearly satisfied. For each small graph H_0 if $\vec{G}' \rightarrow H_0$ then $\vec{G} \rightarrow H_0$ holds again (with the same embedding). Finally since $k > |V(H)|$ and so $\vec{G} \rightarrow \vec{H}$. The same embedding shows us that $\vec{G}' \rightarrow \vec{H}$. Therefore whenever for a small directed graph H_0 we have $\vec{G} \rightarrow H_0$ then we also have $\vec{G}' \rightarrow H_0$ (through the embedding $\vec{G}' \rightarrow \vec{H}$). So property (ii) also holds.

Removing the orientation of each edge of \vec{G}' and keeping just one-one copies of the edges we get the graph G . It trivially satisfies properties (i), (iii) and (iv). To check (ii) remark at first that we already saw that $G \rightarrow H$ therefore direction \leftarrow is done. For the other direction assume that $|V(H_0)| < k$ and $G \rightarrow H_0$. Create \vec{H}_0 from H_0 in the same manner as \vec{H} was made. Then the same embedding $V(G) \rightarrow V(H_0)$ witnesses $\vec{G} \rightarrow \vec{H}_0$. But then $\vec{H} \rightarrow \vec{H}_0$ holds, and the same embedding witnesses $H \rightarrow H_0$. \square

3. THE SPLITTING PROPERTY

In the forthcoming sections we would like to apply some results from [9] to the posets \mathbb{G} and \mathbb{D} to obtain antichains with certain properties related to dualities and partitions of \mathbb{G} and \mathbb{D} . Concerning \mathbb{G} it would be enough just to quote some theorems from [9], but concerning \mathbb{D} we should reformulate them a bit to make them applicable here.

Let $\mathcal{P} = (P, \leq)$ be a poset. [We find it useful sometimes to maintain a distinction between \mathcal{P} and the underlying set P .] We say that a subset $A \subset P$ is *cut-free* in \mathcal{P} provided there are no $y \in A$ and $x, z \in P$ such that $x < y < z$ and $A \cap [x, z] = A \cap ([x, y] \cup [y, z])$. An element $y \in P$ is called *cut-point* iff there are $x, z \in P$ such that $x < y < z$ and $[x, z] = [x, y] \cup [y, z]$. Clearly there is no cut-point in a cut-free set.

If $\mathcal{P} = (P, <)$ is a poset and $A \subset P$ then recall that the *upset* A^\uparrow and the *downset* A^\downarrow of A are the sets

$$A^\uparrow = \{p \in P : \exists a \in A \ a \leq p\}, \quad A^\downarrow = \{p \in P : \exists a \in A \ p \leq a\};$$

also, use this natural extension of the notation,

$$A^\uparrow = A^\uparrow \cup A^\downarrow.$$

As usual, we drop the braces and write a^\uparrow , a^\downarrow , and a^\uparrow in place of $\{a\}^\uparrow$, $\{a\}^\downarrow$ and $\{a\}^\uparrow$, respectively.

A *maximal antichain* A in P is a set of pairwise incomparable elements [an antichain] that is maximal with respect to containment. We say that a maximal antichain A *splits* if there is a partition (B, C) of A such that $P = B^\uparrow \cup C^\downarrow$. We say that A *strongly splits* if and only if there is a partition (B, C) of A such that for each $p \in P \setminus A$ either the set $B \cap p^\downarrow$ or the set $C \cap p^\uparrow$ are infinite.

To construct maximal antichains with desired properties [for instance, splitting or non-splitting], it is useful to be able to extend existing finite antichains to maximal ones in certain special ways. This motivated Erdős and Soukup [9] to formulate this definition: call \mathcal{P} *loose* if for all $x \in P$ and $F \in [P]^{<\omega}$, if $x \notin F^\uparrow$ there is $y \in x^\uparrow \setminus \{x\}$ that is incomparable to all elements in F . This property is the key in showing that very familiar infinite distributive lattices, such as $([\omega]^{<\omega}, \subseteq)$, the lattice of finite subsets of a countably infinite set, do not have the splitting property. We shall see that the nontrivial part of \mathbb{G} has this property [see Theorem 4.1] but that \mathbb{D} requires a sharpening of the definition [see Theorem 5.1].

Definition 3.1. Let $\mathcal{P} = \langle P, \leq \rangle$ be a poset and let $P' \subset P$. We say that P' is a *loose kernel* in \mathcal{P} if

- (LK) for all finite subsets $F \subseteq P'$ and $x \in P \setminus F^\uparrow$ there is $y \in [x^\uparrow \cap P']$, $y \neq x$, such that each element of F is incomparable to y .

Of course, \mathcal{P} is *loose* if P itself is a loose kernel in \mathcal{P} , just as in [9].

Remarks. (1) A loose kernel P' in a poset \mathcal{P} does not have maximal elements – just take $F = \emptyset$ in (LK) in the definition and any x in P' to produce $y \in P'$ such that $y > x$. In particular, P' is infinite. Also, if \mathcal{P} contains a loose kernel then

there is a loose kernel of \mathcal{P} that is maximal, with respect to containment. This is easily shown using Zorn's Lemma.

(2) Regarding the homomorphism poset \mathbb{D} , it is not loose since it has finite maximal antichains – a finite maximal antichain as F in (LK) shows that (LK) fails. Moreover, \mathbb{D} has infinitely many finite maximal antichains, so we cannot obtain a loose kernel for \mathbb{D} by deleting finitely many elements, as we can for \mathbb{G} .

Here is a condition that allows extension of a finite antichain in a particular special way.

Definition 3.2. Let $\mathcal{P} = \langle P, \leq \rangle$ be a poset and $P' \subset P$. We say that P' has the *finite antichain extension property* (in \mathcal{P}) provided

(FAE) for all finite antichains $F \subseteq P'$ and $x \in P \setminus F$ there is $y \in [x^\downarrow \cap P']$ such that each element of F is incomparable to y .

Observation 3.3. If $P' \subset P$ is both a loose kernel in $\mathcal{P} = (P, \leq)$ and a loose kernel in the dual $\mathcal{P}^d = (P, \geq)$ then P' has the finite antichain extension property in \mathcal{P} .

The following observation is a sharpening of [9, Theorem 3.9]. We include the straightforward proof to illustrate how the FAE property can be applied.

Theorem 3.4. Let $\mathcal{P} = \langle P, \leq \rangle$ be a countably infinite poset, let $P' \subset P$ have the finite antichain extension property in \mathcal{P} , and let $A_1 \subset P'$ be a finite antichain. Then there is a strongly splitting \mathcal{P} -maximal antichain A such that $A_1 \subset A \subset P'$.

Proof. Let $\{p_n : n < \omega\}$ be an ω -abundant enumeration of P , that is, the set $\{n : p_n = p\}$ is infinite for each $p \in P$. Let $A_1 = \{a_0, a_1, \dots, a_{r-1}\}$. Proceed by induction on $i \geq r$ to construct an infinite antichain $A = \{a_i : i < \omega\} \subset P'$:

- if $p_i \notin \{a_j : j < i\}$ then let a_i be comparable to p_i ;
- if $p_i \in \{a_j : j < i\}$ then let $n_i = \min\{n : p_n \notin \{a_m : m < i\}\}$ and let a_i be comparable to p_{n_i} .

This construction can be carried out because P' has the finite antichain extension property.

Let $p \in P \setminus A$. Then the set $A_p = \{a_i : p_i = p\}$ is infinite and for each $a \in A_p$ the element a and p are comparable. Let (B, C) be a partition of A such that $|B \cap A_p| = |C \cap A_p| = \omega$ for each $p \in P \setminus A$.

Then the partition (B, C) has the required property. \square

The following results show that the existence of a loose kernel guarantees an infinite non-splitting maximal antichain. The first is a slight generalization of [9, Theorem 3.6].

Theorem 3.5. Let $\mathcal{P} = \langle P, \leq \rangle$ be a countably infinite poset, and let $P' \subset P$ be a loose kernel in \mathcal{P} . Then there exists an infinite non-splitting \mathcal{P} -maximal antichain $A \subset P'$.

Proof. See [9, Theorem 3.6]. \square

Theorem 3.6. *Let $\mathcal{P} = \langle P, \leq \rangle$ be a countably infinite poset, let $P' \subset P$ be a loose kernel in \mathcal{P} , and let $A_1 \subset P'$ be a non-maximal antichain in P . Then there is an infinite non-splitting \mathcal{P} -maximal antichain A such that $A_1 \subset A \subset P'$.*

Proof. The set $P' \setminus A_1^\uparrow$ is a loose kernel in $P \setminus A_1^\uparrow$, and $P \setminus A_1^\uparrow \neq \emptyset$ because A_1 was not a maximal antichain. Hence by Theorem 3.5 there is a $P \setminus A_1^\uparrow$ -maximal antichain $A' \subset P' \setminus A_1^\uparrow$ which does not split in $P \setminus A_1^\uparrow$. Then $A = A_1 \cup A'$ is a maximal antichain in P having the required properties. \square

4. THE HOMOMORPHISM POSET \mathbb{G}

The partially ordered set \mathbb{G} of hom-equivalence classes of finite undirected graphs is known to have only two finite maximal antichains — $\{K_1\}$ and $\{K_2\}$. Consequently, there are no nontrivial dualities. However, in studying the ordered set \mathbb{G} , it is interesting to know whether maximal antichains split and whether antichains extend to maximal ones that split.

Let $\mathbb{G}' = \mathbb{G} \setminus \{K_1, K_2\}$. For any bipartite graph G , $G \rightarrow K_2$, so we know that all graphs in \mathbb{G}' are hom-equivalent to graphs all of whose connected components contain odd cycles. The *odd girth* of a graph G , $\text{oddgirth}(G)$, is the length of the shortest odd cycle contained in the graph. As with girth, if graph does not contain any odd cycles, its oddgirth is regarded as infinite.

The notion of odd girth is useful in dealing with homomorphism questions because of this: for graphs G and H , if $\text{oddgirth}(G) < \text{oddgirth}(H)$ then $G \rightarrow H$. Also, it is straightforward to construct graphs of prescribed odd girth and chromatic number using shift graphs — for instance, see [11, Theorem 2.23]. Alternatively, the original Erdős result could be used in the first part of the proof below.

Theorem 4.1. *\mathbb{G}' is both a loose kernel in \mathbb{G} and a loose kernel in \mathbb{G}^d . Hence, \mathbb{G}' has the finite antichain extension property in \mathbb{G} .*

Proof. Let $\mathcal{F} \subseteq \mathbb{G}'$ be finite. To see that \mathbb{G}' is a loose kernel in \mathbb{G} , let $X \in \mathbb{G} \setminus \mathcal{F}^\uparrow$, that is, $F \rightarrow X$ for all $F \in \mathcal{F}$. Let Y' be a graph such that for all $F \in \mathcal{F}$,

- (i): $\text{oddgirth}(Y') > \text{oddgirth}(F')$ for all components F' of F , and
- (ii): $\chi(Y') > \chi(F)$.

Now let $Y = X + Y'$ and let $F \in \mathcal{F}$. By (i), $F \rightarrow Y$, since no component of F has a homomorphism to Y and $F \rightarrow X$. By (ii), $Y' \rightarrow F$, so $Y \rightarrow F$. Hence, (LK) holds and \mathbb{G}' is loose in \mathbb{G} .

Now let us show that \mathbb{G}' is loose in the dual. Let $H \in \mathbb{G} \setminus \mathcal{F}^\downarrow$, that is, $H \rightarrow F$ for all $F \in \mathcal{F}$. Let $k = \max\{|V(H)|, |V(F)| : F \in \mathcal{F}\}$ and $\ell = \max\{\text{oddgirth}(F) : F \in \mathcal{F}\} + 1$. Here ℓ is finite because $\mathcal{F} \subset \mathbb{G}'$.

By Theorem 2.1 there is a graph $G \in \mathbb{G}$ such that $\text{girth}(G) > \ell$ and for all $K \in \mathbb{G}$ where $|V(K)| \leq k$,

$$G \rightarrow K \iff H \rightarrow K.$$

Therefore $G \rightarrow H$ but for all $F \in \mathcal{F}$ we have $G \not\rightarrow F$. Since $\text{girth}(G) > \text{oddgirth}(F)$ we have $F \not\rightarrow G$ for each $F \in \mathcal{F}$. Furthermore $H \not\rightarrow K_2$ therefore $H \in \mathbb{G}'$,

therefore $G \not\prec K_2$ and so $G \in \mathbb{G}'$. Thus (LK) holds for \mathbb{G}' in \mathbb{G}^d , and we can apply Observation 3.3. \square

As noted above, it is well-known that \mathbb{G}' has no finite maximal antichains. We include a short proof to illustrate the relationship between loose kernels and maximal antichains.

Corollary 4.2. *There is no finite maximal antichain in \mathbb{G}' .*

Proof. Indeed, let $\mathcal{F} \subset \mathbb{G}'$ be a finite antichain. Then $K_1 < F_i$ (for each i) therefore $K_1 \notin \mathcal{F}^\uparrow$. The application of Theorem 4.1 gives us an element of \mathbb{G}' , which is incomparable to \mathcal{F} . \square

Since there are no finite maximal antichains and every finite antichain extends to a maximal one, each finite antichain can be extended to an infinite maximal antichain. The following shows that quite different behaviour can be found in the various extensions.

Corollary 4.3. *Let $A \subseteq \mathbb{G}'$ be a finite antichain. Then*

- (i): *there exists a non-splitting maximal antichain $A_1 \subset \mathbb{G}'$ such that $A \subset A_1$, and*
- (ii): *there exists a strongly splitting maximal antichain $A_2 \subset \mathbb{G}'$ such that $A \subset A_2$.*

Proof. (i): This is a direct consequence of Theorem 3.6, applied to the poset \mathbb{G} and the loose kernel \mathbb{G}' .

(ii): We can apply Theorem 3.4 because \mathbb{G}' has the finite antichain extension property in \mathbb{G} . \square

The notions of a *cut-point* and a *cut-free* subset are closely tied to the splitting property: see [1] and [9]. They have also been studied independently in the context of homomorphism orders of graphs: see [12]. We provide a short proof that \mathbb{G}' is cut-free, both to illustrate an application of the sparse incomparability lemma and to highlight differences between \mathbb{G} and \mathbb{D} that we shall see again in the next section.

Proposition 4.4. *\mathbb{G}' is cut-free.*

Proof. We need to show that for all triples $F < G < H$, if $G \in \mathbb{G}'$ (and therefore $H \in \mathbb{G}'$) then there is a $G' \in \mathbb{G}'$ such that $F < G' < H$ and G' is incomparable to G .

Since $\text{oddgirth}(G)$ is finite for $G \in \mathbb{G}'$, we can apply Theorem 2.1 to H with parameters $k = \max(|V(G)|, |V(F)|) + 1$ and $\ell = \text{oddgirth}(G) + 1$ to get a graph H' such that:

- $H' \rightarrow H$, since $H \rightarrow H$,
- $H' \nrightarrow G$, since $H \nrightarrow G$,
- $H' \nrightarrow F$, since $H \nrightarrow F$, and
- $\text{girth}(H') > \ell$.

Since $\text{oddgirth}(G) < \ell$ we have $X \nrightarrow H'$ for each connected component X of G . Therefore the graph $G' = F + H'$ satisfies the requirements. \square

5. THE HOMOMORPHISM POSET \mathbb{D}

In [10], the complete characterization of finite maximal antichains in the homomorphism poset for finite relational structures with a single relation shows the crucial role of forests. In the study of \mathbb{G} , odd cycles play a crucial role. So, one might hope that the investigation the two subsets \mathbb{D}' and \mathbb{D}^* of \mathbb{D} defined below would lead to the construction of interesting antichains by verifying the loose kernel or FAE properties, then employing results such as Theorems 3.4 – 3.6. It turns out to be a bit more complicated.

Before defining these, it is useful to recall that a finite directed graph X is a *core* if every homomorphism of X to itself is bijective. Every directed graph is homomorphically equivalent to a unique core, and, so, every directed graph class contains exactly one core (cf. [11]). For the rest of this section, we shall use “graph” for “directed graph” and, given a directed graph X , \overline{X} denotes its undirected version.

We now define the two subsets of \mathbb{D} :

- $\mathbb{D}' \subseteq \mathbb{D}$ consists of all graph classes with core X such that every connected component of \overline{X} contains an odd cycle;
- $\mathbb{D}^* \subseteq \mathbb{D}$ consists of all graph classes with core Y such that at least one connected component of \overline{Y} has contains an odd cycle.

Of course, $\mathbb{D}' \subseteq \mathbb{D}^*$, while the graph H defined below in the proof of Proposition 5.2 is in \mathbb{D}^* and not in \mathbb{D}' .

The following result collects some straightforward observations about these subsets in \mathbb{D} . The proofs have been omitted since the methods are not very different from those encountered in the undirected case.

Proposition 5.1. *In the partially ordered set \mathbb{D} :*

- (i): \mathbb{D}' is a loose kernel in \mathbb{D} ;
- (ii): \mathbb{D}^* is a cut-free subset of \mathbb{D} ; and, *MISTAKEN*
- (iii): \mathbb{D}^* is loose in $(\mathbb{D}^*)^d$.

ADDENDUM

Lemma F *Let $\mathcal{F} \in [\mathbb{D}']^{<\omega}$ and let $X \in \mathbb{D} \setminus \mathcal{F}^\uparrow$. Then there is an $X' \in \mathbb{D}'$ such that $X \rightarrow X'$ and $F \not\rightarrow X'$ for each $F \in \mathcal{F}$.*

Proof. If X has components without odd cycles, then glue one-one big odd cycles to one-one vertices of those components. Then, of course, $X' > X$. If these cycles has length $> |V(F)|$ for every $F \in \mathcal{F}$ then $F \not\rightarrow X'$ for every $F \in \mathcal{F}$. \square

Theorem G \mathbb{D}' is an upwards loose kernel in \mathbb{D} .

Proof. Let $\mathcal{F} \in [\mathbb{D}']^{<\omega}$ and let $X \in \mathbb{D} \setminus \mathcal{F}^\uparrow$. We are looking for a $Y \in \mathbb{D}'$ such that $X \rightarrow Y \not\rightarrow X$ and Y is incomparable to each $F \in \mathcal{F}$. If $X \notin \mathbb{D}'$ then take X' from Lemma F. If we can construct the required Y for X' then it works with X as well. Therefore we can assume that each component of X contains odd (undirected) cycle. Let $m := \max\{|F| : F \in \mathcal{F}\} + 1$ and $\ell := \max\{\text{oddgirth}(F) : F \in \mathcal{F}\} + 1$. Apply Theorem 2.2 for an arbitrary orientation of the complete graph K_m with these parameters. The result is a directed connected graph G with a finite oddgirth $> \ell$, such that $G \not\rightarrow F \in \mathcal{F}$ (since $K_m \not\rightarrow F$, due its chromatic number $\chi(F) < m$). Furthermore each component in F has odd cycle therefore $F \not\rightarrow G$ as well. Finally the disjoint union $Y = G + X$ satisfies the required conditions. \square

Nešetřil and Tardiff, [17], described all maximal finite antichains in the direct homomorphism poset, therefore the following statement is a well-known fact:

Corollary H *There is no finite maximal antichain in \mathbb{D}' .*

Proof. Like Corollary 4.2. \square

Corollary I *Let $A_1 \in [\mathbb{D}']^{<\omega}$ be a non-maximal antichain in \mathbb{D} . Then there is a maximal antichain $A \subset \mathbb{D}'$ such that $A_1 \subset A$ and A does not split, moreover $(A \setminus A_1) \subset B$ and C is infinite for all $\langle B, C \rangle \in \mathfrak{S}(A)$.*

Proof. Due to Theorem G this is just the direct application of Theorem 3.5 for the poset \mathbb{D} and the upwards loose kernel \mathbb{D}' . \square

We have seen so far that \mathbb{D}' is upward loose kernel in \mathbb{D} , however we can not prove that \mathbb{D}' is a downward loose kernel in \mathbb{D} , or that \mathbb{D}' is cut-free in \mathbb{D} .

On the other hand, as we will seen soon, \mathbb{D}^* is cut-free in \mathbb{D} , \mathbb{D}^* is "almost" a downward loose kernel in \mathbb{D} , but we will not know if \mathbb{D}^* is upward loose kernel in \mathbb{D} or not.

Theorem J *\mathbb{D}^* is cut-free in \mathbb{D} .*

Proof. We should show that for all triple $F < G < H$ if $G \in \mathbb{D}^*$ (and therefore $H \in \mathbb{D}^*$ as well) there is a $G' \in \mathbb{D}^*$ such that $F < G' < H$ and G' is incomparable to G .

Apply Theorem 2.2 for the graph H with parameters $m = |V(G)| + 1$ and $\ell = \text{oddgirth}(G) + 1$. Since $H \not\rightarrow G$ therefore $H' \not\rightarrow G$ as well. Due to the big odd girth $G \not\rightarrow H'$ also holds (as it shows by the component of G which has an odd cycle). Therefore the graph $G' = F + H'$ satisfies the requirements. \square THIS IS MISTAKEN.

Theorem K *Let $\mathcal{F} \in [\mathbb{D}^*]^{<\omega}$ and $X \in \mathbb{D}$ such that $X \notin \mathcal{F}^\downarrow$. Then there exists $X' \in \mathbb{D}$ such that $X' < X$ and X' is incomparable to each $F \in \mathcal{F}$. Furthermore if $X \in \mathbb{D}^*$ then $X' \in \mathbb{D}^*$ holds as well.*

Proof. Apply Theorem 2.2 for the graph X with parameters $m = \max\{|V(F)| : F \in \mathcal{F}\} + 1$ and $\ell = \max\{\text{oddgirth}(F) : F \in \mathcal{F}\} + 1$. Then $X' \not\rightarrow F \in \mathcal{F}$ since $X \not\rightarrow F$. Furthermore if $F \notin \mathbb{D}^*$ then the same applies for X' therefore $F \not\rightarrow X'$ since F has odd cycle. However, if X has an odd cycle then all odd cycles of X' (and there are such objects) are longer than $\text{oddgirth}(F)$ therefore $F \not\rightarrow X'$ holds again. \square

Unfortunately, we can also prove that

Proposition 5.2. *In the partially ordered set \mathbb{D} :*

- (i): \mathbb{D}' does not have the finite antichain extension property in \mathbb{D} ; and
- (ii): \mathbb{D}^* does not have the finite antichain extension property in \mathbb{D} .

Proof. (i). Let \vec{T}_3 be the transitive tournament on three vertices and let $\vec{P}_3 = (W, F)$ be the directed path on four vertices: $W = \{x_0, x_1, x_2, x_3\}$ and $F = \{(x_i, x_{i+1}) : i = 0, 1, 2\}$. Now let $H = \vec{T}_3 + \vec{P}_3$, the disjoint union of T and P_3 . Then H is a core in \mathbb{D} . Also, $\vec{T}_3 \in \mathbb{D}'$, is a core and $\vec{T}_3 < H$. Regard $\{\vec{T}_3\}$ as an antichain. If \mathbb{D}' had the (FAE) property there would exist a core $H' \in \mathbb{D}'$ such that H' is incomparable to \vec{T}_3 and $H' < H$. However, every connected component of $\overline{H'}$ contains an odd cycle, so $H' < H$ implies that $H' \leq \vec{T}_3$ since no component of H' can be mapped by a homomorphism into \vec{P}_3 .

(ii) We can base an example on any oriented tree T but to be a bit more specific let $T = \vec{P}_k$ be a directed path on k vertices. Then the dual $D(\vec{P}_k)$ is the transitive tournament \vec{T}_k on k vertices (see, for instance [11, Proposition 1.20]). Let $H = \vec{T}_k + \vec{P}_k$. As long as $k \geq 3$, $H \in \mathbb{D}^*$. Also, H and \vec{T}_k are cores in \mathbb{D} and $\vec{T}_k < H$. Regard $\{H\}$ as an antichain. If \mathbb{D}^* had the (FAE) property there would exist a core $H' \in \mathbb{D}^*$ such that H' is incomparable to H and $\vec{T}_k < H'$. But (\vec{P}_k, \vec{T}_k) is a dual pair, so the fact that H' is not below \vec{T}_k implies $\vec{P}_k < H'$ – just apply equation (3) in this special case. From this it would follow that $H = \vec{T}_k + \vec{P}_k < H'$, a contradiction. \square

Fortunately there is another subset \mathbb{D}^c of \mathbb{D} which is both an upward loose kernel in \mathbb{D} and has the finite antichain extension property in \mathbb{D} . To discuss it, first we need an easy observation. A finite directed graph \vec{C} is a *directed cycle* if it is connected and each vertex has indegree and outdegree 1. It is easily seen that each directed cycle is a core. Use \vec{C}_k to denote the directed cycle on k vertices.

Proposition 5.3. *Let \vec{C} be a directed cycle and T be a graph such that \overline{T} is an arbitrary tree. Then $T \rightarrow \vec{C}$.*

Proof. Map a vertex v of T to any vertex of the cycle. Next the vertices adjacent to v in \overline{T} can be mapped into vertices of \vec{C} so that directed edges are preserved. Since there is no cycle in \overline{T} we can finish the process easily. \square

Let \mathbb{D}^c be the set of all homomorphism classes in \mathbb{D} whose core X has the property that for some \vec{C} , $\vec{C} \rightarrow X$. Here is a direct consequence of Proposition 5.3.

Observation 5.4. *Let $G \in \mathbb{D}^c$ and let $T \in \mathbb{D}$ be an oriented tree. Then $G + T \sim G$.*

Hence we can assume that no component of an element of \mathbb{D}^c can be embedded into a tree. Therefore from now on we assume that each component X of each element of \mathbb{D}^c has the property that \overline{X} contains a cycle.

Theorem 5.5. *\mathbb{D}^c is a loose kernel in \mathbb{D} .*

Proof. Let $\mathcal{F} \subseteq \mathbb{D}^c$ be finite, and $X \in \mathbb{D}$ but $X \notin \mathcal{F}^\uparrow$. We are going to find an $Y \in \mathbb{D}^c$ such that $X \rightarrow Y$, $Y \not\rightarrow X$, and Y is incomparable to each $F \in \mathcal{F}$.

Let $n := \max\{|X|, |F| : F \in \mathcal{F}\}$. Using the Erdős theorem, obtain a graph Z such that $\chi(\overline{Z}) > n$, $\text{girth}(\overline{Z}) > n$, \overline{Z} is connected, and Z contains at least one directed cycle. Then $Z \in \mathbb{D}^c$. Let $Y = X + Z$. Since $Z \in \mathbb{D}^c$ therefore $Y \in \mathbb{D}^c$ as well. Clearly $X \rightarrow Y$ while $Y \not\rightarrow X$ because $\chi(Y) > |X|$.

Assume that f is a homomorphism of F to Y . Then there is a component K of F such that $f[K] \subseteq Z$. But $|V(K)| \leq n$ while $\text{girth}(Z) > n$, hence the image $f[K]$ is a tree, which contradicts the assumption that no component of an element of \mathbb{D}^c can be mapped into a tree. \square

Theorem 5.6. *Let $A_1 \subseteq \mathbb{D}^c$ be a finite antichain. Then there is a non-splitting antichain A such that $A_1 \subseteq A \subseteq \mathbb{D}^c$ and A is maximal in \mathbb{D} .*

Proof. Since \mathbb{D}^c is an upward loose kernel it can be used in Theorem 3.5 to extend a non-maximal antichain into a non-splitting antichain, maximal in \mathbb{D} . \square

Theorem 5.7. *\mathbb{D}^c has the finite antichain extension property in \mathbb{D} .*

Proof. Let $\mathcal{F} \subseteq \mathbb{D}^c$ be a finite antichain and $X \in \mathbb{D}$. We need to find $Y \in (X^\uparrow \cap \mathbb{D}^c) \setminus \mathcal{F}^\uparrow$. In case $X \notin \mathcal{F}^\uparrow$ then Theorem 5.5 provides the required Y .

Assume now that $X \in \mathcal{F}^\uparrow$. Then $X \in \mathbb{D}^c$ because there exists $F \in \mathcal{F}$ with $F \rightarrow X$ and the image of its directed cycle of F is a directed cycle in X . Let us assume that X contains the directed cycle \vec{C}_k .

Let $n = \max\{|X|, |F| : F \in \mathcal{F}\}$. Apply Theorem 2.2 with $H = X$ and $m = \ell = n$ to obtain $X' = H'$. Now let $Y = X' + \vec{C}_{kn}$. Then $X' \rightarrow X$ and $\vec{C}_{kn} \rightarrow \vec{C}_k$ therefore $Y \rightarrow X$. At the same time $X \not\rightarrow Y$ since $\text{girth}(Y) > \ell \geq |X|$. Thus, the cycle \vec{C}_k of X cannot be embedded into Y . The same applies for the directed cycles in each $F \in \mathcal{F}$. Therefore, $F \not\rightarrow Y$. Finally we have $X \not\rightarrow F$ and so $Y \not\rightarrow F$. \square

Corollary 5.8. *\mathbb{D}^c does not contain finite maximal antichains.*

We recall that a full description of the finite maximal antichains in \mathbb{D} is given in [10].

Corollary 5.9. *Let $A_1 \subseteq \mathbb{D}^c$ be a finite antichain in \mathbb{D} . Then there is a strongly splitting \mathbb{D} -maximal antichain $A_1 \subset A \subset \mathbb{D}^c$.*

Proof. This is just the direct application of Theorem 3.4 to the posets \mathbb{D} and \mathbb{D}^c . \square

Here is a final use of our methods in describing the order structure of \mathbb{D} .

Theorem 5.10. *\mathbb{D}^c is cut-free in \mathbb{D} .*

Proof. Let $F < G < H$ where $G \in \mathbb{D}^c$ (and therefore $H \in \mathbb{D}^c$ as well). We need a $G' \in \mathbb{D}^c$, which is incomparable to G but $F < G' < H$.

Let $n = \max\{|F|, |G|, |H|\}$. Apply Theorem 2.2 to H with parameters $m = \ell = n$ to obtain the directed graph H' . Since $H \in \mathbb{D}^c$, there is k such that \vec{C}_k is a subgraph of H . Let $G' = F + H' + \vec{C}_{kn}$.

Then $F \rightarrow G'$ since F is a subgraph of G' . Furthermore $H' \rightarrow H$ due to the fact that $|H| = n \leq m$ and $H \rightarrow H$. Due to our assumption on \mathbb{D}^c , each component of the graph G contains cycles, and at least one of them, say K , cannot be embedded into F . Therefore if $G \rightarrow Y$ then for this component we have $K \rightarrow H' + C_{nk}$. However, $\text{girth}(H' + C_{nk}) > |K|$, hence K is embedded into a tree, a contradiction. \square

REFERENCES

- [1] R. Ahlswede, P.L. Erdős, N. Graham, A splitting property of maximal antichains, *Combinatorica* **15** (1995) 475–480
- [2] D. Duffus, B. Sands, Minimum sized fibres in distributive lattices, *J. Austral. Math. Soc.* **70** (2001) 337–350
- [3] D. Duffus, B. Sands, Splitting numbers of grids, *Electronic J. Comb.* **12** (2005) #R17
- [4] M. Džamonja, A note on the splitting property in strongly dense posets of size \aleph_0 , *Radovi Matematički* **8** (1999) 321–326
- [5] Paul Erdős: Graph theory and probability. *Canad. J. Math.* **11** (1959) 34–38
- [6] Paul Erdős, A. Hajnal: On chromatic number of graphs and set systems, *Acta Math. Hungar.* **17** (1966), 61–99
- [7] Paul Erdős, A. Rényi: On random graphs I. *Publ. Math. Debrecen* **6** (1959) 290–297
- [8] P.L. Erdős, Splitting property in infinite posets, *Discrete Math.* **163** (1997) 251–256
- [9] P.L. Erdős, L. Soukup: How to split antichains in infinite posets, to appear in *Combinatorica* **27** (2) (2007).
- [10] J. Foniok, J. Nešetřil, C. Tardif: Generalized dualities and maximal finite antichains in the homomorphism order of relational structures, to appear in *European J. Comb.*; extended abstract in *Lecture Notes in Computer Science* **4271** (2006) 27 – 36.
- [11] P. Hell, J. Nešetřil, *Graphs and Homomorphisms*, Oxford University Press, Oxford (2004)
- [12] J. Nešetřil, Sparse incomparability for relational structures, in preparation
- [13] J. Nešetřil and P. Ossona de Mendez, Cuts and bounds, *Discrete Math.* **302** 1 – 3 (2005) 211 – 224
- [14] J. Nešetřil, A. Pultr, On classes of relations and graphs determined by subobjects and factorobjects, *Discrete Math.* **22** (1978) 287–300
- [15] J. Nešetřil, V. Rödl: Chromatically optimal rigid graphs, *J. Comb. Theory Ser. (B)* **46** (1989) 133–141
- [16] J. Nešetřil, C. Tardif: Duality theorems for finite structures (characterising gaps and good characterisations) *J. Comb. Theory Ser. (B)* **80** (2000) 80–97
- [17] J. Nešetřil, C. Tardif: On maximal finite antichains in the homomorphism order of directed graphs. *Discuss. Math. Graph Theory* **23** (2003) 325–332
- [18] J. Nešetřil, Xuding Zhu: On sparse graphs with given colorings and homomorphisms, *J. Comb. Theory Ser. (B)* **90** (2004), 161–172
- [19] A. Pultr, V. Trnková: *Combinatorial, Algebraic and Topological Representations of Groups, Semigroups and Categories*, North-Holland, Amsterdam (1980)
- [20] E. Welzl: Color families are dense, *Theoret. Comput. Sci.* **17** (1982) 29–41

(Dwight Duffus) MATHEMATICS & COMPUTER SCIENCE DEPARTMENT, EMORY UNIVERSITY,
ATLANTA, GA USA, 30322

E-mail address, Dwight Duffus: `dwight@mathcs.emory.edu`

(Péter L. Erdős and Lajos Soukup) ALFRÉD RÉNYI INSTITUTE OF MATHEMATICS, BUDAPEST,
P.O. BOX 127, H-1364 HUNGARY

E-mail address, Péter L. Erdős and Lajos Soukup: `elp@renyi.hu, soukup@renyi.hu`

(Jaroslav Nešetřil) APPLIED MATHEMATICS MATHEMATICS DEPARTMENT & INSTITUTE OF
THEORETICAL COMPUTER SCIENCE, CHARLES UNIVERSITY, PRAGUE, CZECH REPUBLIC

E-mail address, Jaroslav Nešetřil: `nesetril@kam.mff.cuni.cz`