

Optimal Embedded and Enclosing Isosceles Triangles

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Given a triangle Δ , we study the problem of determining the smallest enclosing and largest embedded isosceles triangles of Δ with respect to area and perimeter. This problem was initially posed by Nandakumar [17, 22] and was first studied by Kiss, Pach, and Somlai [13], who showed that if Δ' is the smallest area isosceles triangle containing Δ , then Δ' and Δ share a side and an angle. In the present paper, we prove that for any triangle Δ , every maximum area isosceles triangle embedded in Δ and every maximum perimeter isosceles triangle embedded in Δ shares a side and an angle with Δ . Somewhat surprisingly, the case of minimum perimeter enclosing triangles is different: there are infinite families of triangles Δ whose minimum perimeter isosceles containers do not share a side and an angle with Δ .

Keywords: Isosceles triangle; special container; minimal cover.

1. Introduction

The following classical problem is the starting point of our investigation. Given two convex bodies, C and C' in \mathbb{R}^d , decide whether C can be moved into a position

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where it covers C'. One can easily list some necessary conditions, for instance, the volume, the surface area and the diameter of C has to be at least as large as the one of C'. However, solving the decision problem can be rather challenging, even in \mathbb{R}^2 , or for special cases that might seem friendly at first sight.

For instance, consider the setup where C' is the 'shadow' of C, that is, C is embedded into \mathbb{R}^3 and C' is the orthogonal projection of C onto a 2-dimensional affine subspace. The necessary conditions are clearly satisfied and it looks plausible that there is always a congruent copy of C which covers C'. However, the proof of this fact is far from straightforward [5, 14], and curiously, the result does not generalize to higher dimensions: for $d \geq 3$, no convex d-polytope embedded in \mathbb{R}^{d+1} can cover all of its shadows [5].

Another special case is where both convex bodies are triangles in \mathbb{R}^2 : given two triangles Δ and Δ' , the goal is to find an efficient way to decide whether Δ can be brought into a position where it covers Δ' . This is a classical problem posed by Steinhaus [26] in 1964 and an algorithmic solution was proposed only 29 years later by Post [21], who described a set of 18 polynomial inequalities of degree 4 such that a copy of Δ can cover Δ' if and only if at least one of these inequalities are satisfied. The key geometric component of Post's solution is the following.

Lemma 1 (Post [21]). If a triangle Δ can be moved to a position where it covers another triangle Δ' , then one can also find a covering position of Δ with a side that contains one side of Δ' .

Results of this kind help us to reduce the number of configurations to consider, and are of both theoretical and practical interest.

1.1. Optimal covers from a class

In the present paper, we study a variant of the covering problem where the body C (or C') is not fixed, but can be chosen from a family of possible objects and we want to find a solution which is in some sense optimal, for example, has minimum area or perimeter.

Several classical problems in geometry can be viewed as covering problems of this kind: finding an optimal enclosing triangle, polygon, or ellipse (Löwner-John ellipse) for a given input set [3, 4, 6–8, 10, 11, 23, 24] as well as their higher dimensional analogues (that is, simplices, polytopes, ellipsoids [12, 19, 28]). Apart from their theoretical interest, these problems have found applications in various areas of computer science and mathematics (optimization, packing and covering, approximation algorithms, convexity, computational geometry), see [9, 15, 25]. In the past decade, several explicit algorithms were proposed for the case of triangles [3, 16, 20, 27].

In this work, we consider containers that are in some sense *oriented*. Apart from ellipses, isosceles triangles are perhaps the most natural candidates for such containers: their orientation is determined by their axis of symmetry. The study of isosceles containers was initiated by Nandakumar [17, 18, 22] who raised the following two optimisation problems:

(A) Given a triangle Δ , determine the minimum area and the minimum perimeter isosceles triangles that contain Δ .

In what follows, we study these two questions, together with their 'dual' versions:

(B) given a triangle Δ , determine the maximum area and the maximum perimeter isosceles triangles embedded (that is, contained) in Δ .

Our goal is to describe the list of possible solutions for each of the above four problems. There are 9 natural candidates for isosceles enclosing or embedded triangles which we call *special* (for a discussion of cases and figures, see Sec. 2).

Definition 2. For a triangle Δ , we say that Δ' is a special enclosing (or embedded) isosceles triangle of Δ if Δ' is isosceles, it encloses (or embeds into) Δ , and it shares a side with Δ and an angle at one end of this side.

Minimum area isosceles containers have been recently studied by Kiss, Pach, and Somlai [13]. They showed that in this case, any optimal container is special and, for each Δ , there are only 3 special isosceles containers for which the minimum can be attained. In this paper, we complete the picture: we characterize the optimal solutions for the other three problems stated above.

We prove that for three of the four optimisation problems considered, the optimum is always attained at a special configuration.

Theorem 3. Let Δ be a triangle in \mathbb{R}^2

- (i) If Δ" ⊇ Δ is a minimum area isosceles container of Δ, then Δ" is a special container of Δ [13].
- (ii) If Δ' ⊆ Δ is a maximum area embedded isosceles triangle in Δ, then Δ' is a special embedded isosceles triangle of Δ.
- (iii) If Δ' ⊆ Δ is a maximum perimeter embedded isosceles triangle in Δ, then Δ' is a special embedded isosceles triangle of Δ.

Moreover, in each of these cases, we will restrict the number of possible optimal configurations to only 3 special triangles (see Corollary 8 and Remarks 9, 12).

Interestingly, in the case of minimum perimeter containers, the optimal triangle is not necessarily special.

Theorem 4. There are infinite families of triangles Δ such that none of their minimum perimeter isosceles containers is special. We describe 5 different types of isosceles containers such that any triangle Δ has a minimum perimeter isosceles container Δ' which belongs to one of these types. Only 3 out of them are special.

We note that for any input triangle, the areas and perimeters of the at most 5 possible optimal embedded or enclosing isosceles triangles can be computed efficiently.

Our paper is organized as follows. In Sec. 2, we fix the notation and list some easy preliminary statements. In Sec. 3 and Sec. 4, we present the proofs of Theorem 3(ii) and Theorem 3(iii), respectively. Finally, Sec. 5 is dedicated to the description of the 5 types of isosceles containers mentioned in Theorem 4 and to the proof of this result.

2. Preliminaries and Notation

We start by stating three simple lemmas (Lemmas 5–7) formulating elementary properties of the optimal embedded and enclosing isosceles triangles and, thus, providing a unified starting point for the proofs of Theorems 3 and 4. The straightforward proofs of these lemmas can be found in the appendix to the arXiv version of the paper; see also [1].

Lemma 5. Let Δ_1 and Δ_2 be two triangles.

- (i) Let Δ'₁ be a similar copy of Δ₁ of maximum area (resp. perimeter) such that Δ'₁ ⊆ Δ₂. Then
 - (a) there is a side of Δ_2 that contains a side of Δ'_1 ;
 - (b) every side of Δ_2 contains a vertex of Δ'_1 ;
 - (c) Δ'_1 and Δ_2 have a common vertex.
- (ii) Let Δ'₁ be a similar copy of Δ₁ of minimum area (resp. perimeter) such that Δ₂ ⊆ Δ'₁. Then
 - (a) there is a side of Δ'_1 that contains two vertices of Δ_2 ;
 - (b) every side of Δ'_1 contains a vertex of Δ_2 ;
 - (c) Δ'_1 and Δ_2 have a common vertex.

Optimal isosceles enclosing and embedded triangles satisfy further properties.

- **Lemma 6.** (i) For every triangle Δ , there exist a minimum area (resp. perimeter) isosceles container of Δ and a maximum area (resp. perimeter) isosceles triangle embedded in Δ .
- (ii) If Δ₁ is a maximum area (resp. perimeter) isosceles triangle embedded in Δ, then every vertex of Δ₁ lies on a side of Δ.
- (iii) If Δ₂ is a minimum area (resp. perimeter) isosceles container of Δ, then every vertex of Δ lies on a side of Δ₂.

Basic notation. For any two points, A and B, let AB denote the closed segment connecting them, and let |AB| stand for the length of AB. To unify the presentation, in the sequel we fix a triangle ABC with side lengths a = |BC|, b = |AC|, c = |AB|. If two sides are of the same length, then ABC is the unique minimum area and

perimeter isosceles container (and also maximum area and perimeter embedded isosceles triangle) of itself. Therefore without loss of generality, we assume that a < b < c. In the remaining part of this section, we introduce the notation that we use for the special embedded and enclosing isosceles triangles, which are the key objects of our paper (see Definition 2, Theorems 3 and 4).

2.1. Special embedded isosceles triangles

Given a triangle ABC, we describe its special embedded isosceles triangles, that is, all those isosceles triangles contained in ABC that have a common side with ABC and share an angle at one of the endpoints of the common side.

Special embedded triangles of the first kind. Let A' be a point of AC with |A'C| = |BC| and let B' and A'' be two points of AB such that |AB'| = |AC| and |A''B| = |BC| (see Fig. 1). We say that A'BC, AB'C, and A''BC are the special embedded triangles of the first kind associated with ABC.

Special embedded triangles of the second kind. Let C_1 be the intersection of the perpendicular bisector of AB and the segment AC. Analogously, let A_1 be the intersection of the perpendicular bisector of BC and AC, and let B_1 be the intersection of the perpendicular bisector of BC and the line AC (see Fig. 2). The triangles A_1BC , AB_1C , and ABC_1 are the special embedded triangles of the second kind associated with ABC.

Special embedded triangles of the third kind. Let \overline{A} be a point of AB, where $|\overline{A}C| = |BC|$. Analogously, let $\overline{\overline{A}} \in AC$, and $\overline{B} \in BC$ such that $|\overline{\overline{A}}B| = |BC|$, and $|\overline{B}A| = |AC|$ (see Fig. 3). Note that if ABC is non-acute, then $\overline{\overline{A}BC}$ and $A\overline{BC}$

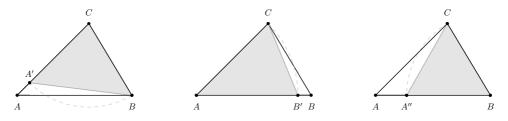


Fig. 1. Special embedded triangles of the first kind.

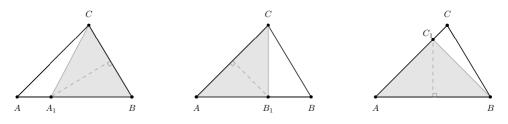


Fig. 2. Special embedded triangles of the second kind.

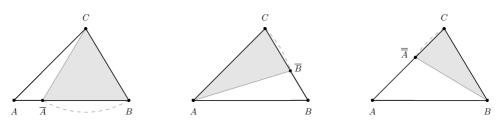


Fig. 3. Special embedded triangles of the third kind.

do not exist. The triangles \overline{ABC} , $\overline{\overline{A}BC}$, and $A\overline{BC}$ are called the *special embedded* triangles of the third kind associated with ABC.

2.1.1. Basic inequalities for special embedded triangles

We collect a few inequalities on the area and perimeter of special isosceles embedded triangles. For a triangle Δ , let $per(\Delta)$ and $area(\Delta)$ denote the perimeter and the area of Δ , respectively.

Lemma 7. If ABC satisfies a < b < c, then

- (i) $\operatorname{area}(A''BC) < \operatorname{area}(A'BC);$
- (ii) $\operatorname{area}(A_1BC) < \operatorname{area}(AB'C)$ and $\operatorname{area}(AB_1C) < \operatorname{area}(ABC_1)$;
- (iii) $\operatorname{area}(\overline{ABC}) < \operatorname{area}(ABC_1)$, $\operatorname{area}(\overline{\overline{ABC}}) < \operatorname{area}(A\overline{\overline{BC}})$, and $\operatorname{area}(A\overline{\overline{BC}}) < \operatorname{area}(AB'C)$;
- (iv) if ABC is obtuse, then $\operatorname{area}(A'BC) < \operatorname{area}(ABC_1)$.

Lemma 7 implies that only 3 of the special embedded triangles of ABC can be optimal.

Corollary 8. If ABC satisfies a < b < c, then any maximum area special embedded triangle of ABC is one of the following triangles: A'BC, AB'C, ABC_1 .

Remark 9. Similar results hold for the perimeter function, implying that any maximum perimeter special embedded triangle of ABC is one of the triangles AB'C, A_1BC , or ABC_1 .

2.2. Special enclosing isosceles triangles

Given a triangle ABC, now we describe its special enclosing isosceles triangles, that is, all those isosceles triangles containing ABC that have a common side with ABCand share an angle at one of the endpoints of the common side.

Special containers of the first kind. Let B' denote the point on the ray \vec{CB} , for which |B'C| = |AC|. Analogously, let C' (and C'') denote the points on \vec{AC} (resp., \vec{BC}) such that |AC'| = |AB| (resp., |BC''| = |AB|), see Fig. 4. We call the

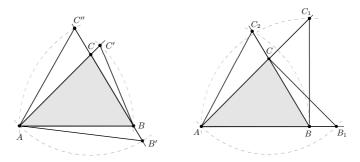


Fig. 4. Special containers of the first kind (AB'C, ABC', and ABC'') and second kind $(AB_1C, ABC_1, \text{ and } ABC_2)$.

triangles AB'C, ABC', and ABC'' special containers of the first kind associated with ABC.

Special containers of the second kind. Let B_1 denote the point on the ray AB, different from A, for which $|B_1C| = |AC|$. Analogously, let C_1 (resp., C_2) denote the point on \vec{AC} (resp., \vec{BC}) for which $|BC_1| = |AB|$ and $C_1 \neq A$ (resp., $|AC_2| = |AB|$ and $C_2 \neq B$), see Fig. 4. The triangles AB_1C , ABC_1 , and ABC_2 are called the special containers of the second kind associated with ABC.

Special containers of the third kind. Let A be the intersection of the perpendicular bisector of BC and the line AC. Since we have b = |AC| < |AB| = c, the point \overline{A} lies outside of ABC. Analogously, denote by \overline{B} (resp., \overline{C}) the intersection of the perpendicular bisector of AC (resp. AB) and the line BC. (If ABC is non-acute \overline{ABC} and $A\overline{BC}$ do not contain ABC (Fig. 5).) The triangles \overline{ABC} , $A\overline{BC}$, and $AB\overline{C}$ are called the *special containers of the third kind* associated with ABC, provided that they contain ABC.

2.2.1. Basic inequalities for special containers

Similarly to the case of maximum area embedded triangles, we can show that not all special containers can be of minimum perimeter.

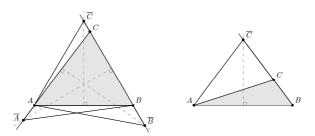


Fig. 5. Special containers of the third kind $(\overline{ABC}, A\overline{BC}, AB\overline{C})$ in the acute and in the non-acute cases.

Lemma 10. If ABC satisfies a < b < c, then

- (i) $\operatorname{per}(ABC') < \operatorname{per}(ABC'')$ and $\operatorname{per}(AB'C) < \operatorname{per}(AB_1C)$;
- (ii) $\operatorname{per}(ABC') < \operatorname{per}(ABC_2) < \operatorname{per}(ABC_1);$
- (iii) $\operatorname{per}(ABC') < \operatorname{per}(\overline{ABC}) < \operatorname{per}(A\overline{B}C)$.

The straightforward proof of Lemma 10 is given in the arXiv version of the paper [1]. Lemma 10 immediately gives the following corollary.

Corollary 11. If ABC satisfies a < b < c, then any minimum perimeter special container of ABC is one of the following triangles: AB'C, ABC', $AB\overline{C}$.

Remark 12. Similar results hold for the area function, implying that a minimum area special container of ABC is one of the triangles AB'C, ABC', or AB_1C .

3. Maximum Area Embedded Isosceles Triangles — Proof of Theorem 3(ii)

Let ABC be a triangle and let XYZ denote one of its maximum area isosceles embedded triangles. In this section, we prove that XYZ has to be a special embedded triangle. We use the notation a = |BC|, b = |AC|, c = |AB|, x = |YZ|, y = |XZ|, z = |XY|, and assume (with no loss of generality) that a < b < c.

By Lemmas 5 and 6, we have the following statements on maximum area embedded isosceles triangles.

Lemma 13. Let XYZ be any maximum area isosceles triangle embedded in ABC. Then

- (i) a side of ABC contains a side of XYZ;
- (ii) every side of ABC contains a vertex of XYZ;
- (iii) ABC and XYZ have a common vertex;
- (iv) no vertex of XYZ lies in the interior of ABC.

If XYZ has at least two common vertices with ABC, then by Lemma 13(iv), XYZ and ABC have a common side and a common angle. Therefore, we can assume that ABC and XYZ have exactly one common vertex.

Denote the midpoints of the sides BC, AC, and AB by m_A , m_B , and m_C , respectively. We divide the boundary of ABC into 3 polylines defined as

$$\widehat{m_A m_B} = m_A C \cup C m_B, \quad \widehat{m_B m_C} = m_B A \cup A m_C, \quad \widehat{m_C m_A} = m_C B \cup B m_A.$$

We get the following constraint on the position of X, Y, and Z:

Lemma 14. Let XYZ be a maximum area embedded isosceles triangle of the triangle ABC. Then each of $\widehat{m_A m_B}, \widehat{m_B m_C}$, and $\widehat{m_C m_A}$ contains exactly one vertex of XYZ.

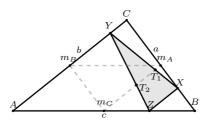


Fig. 6. Illustration for the proof of Lemma 14.

Proof. By Lemma 13, X, Y, Z lies on the boundary of *ABC*. Assume, without loss of generality, that $\widehat{m_A m_C}$ contains X and Z, see Fig. 6.

Let $T_1 = m_A m_C \cap XY$ and $T_2 = m_A m_C \cap YZ$. Then $\operatorname{area}(XT_1T_2Z) \leq \operatorname{area}(Bm_Am_C)$ and by $|T_1T_2| \leq |m_Am_C|$ we obtain that $\operatorname{area}(T_2T_1Y) \leq \operatorname{area}(m_Am_Bm_C)$. Thus we have

$$\operatorname{area}(XYZ) \le \operatorname{area}(Bm_Am_C) + \operatorname{area}(m_Am_Bm_B) = \frac{\operatorname{area}(ABC)}{2}$$

On the other hand, since $c \leq a + b \leq 2b$, the special embedded triangle AB'C satisfies

$$\operatorname{area}(AB'C) = \frac{b^2 \sin(\triangleleft CAB)}{2} > \frac{bc \sin(\triangleleft CAB)}{4} = \frac{\operatorname{area}(ABC)}{2}.$$

Hence, $\operatorname{area}(XYZ) < \operatorname{area}(AB'C)$, which contradicts the maximality of the area of XYZ.

Lemmas 13 and 14 imply that a maximum area embedded isosceles triangle of ABC is either special or its vertex arrangement corresponds to one of the 9 cases illustrated in Fig. 7.

To complete the proof of Theorem 3(ii), it remains to prove that none of the arrangements depicted on Fig. 7 can be optimal. We prove this for each of the 9 cases, separately. Note that in some instances, we will refer to special embedded triangles using their specific labeling introduced in Sec. 2.1.

Case A: The common vertex of ABC and XYZ is A = X.

Subcase A.1: $Y \in BC$ and $Z \in AC$.

Observe that since b < c, the orthogonal projection of A onto CB is contained in Cm_A , which implies that $\langle AYB \rangle$ is obtuse. Thus, we can rotate XYZ about X such that two of its vertices get to the interior of ABC and so, by Lemma 13, XYZ cannot be of maximum area.

Subcase A.2: Both Y and Z are in BC.

If y = z, then we can increase area(XYZ) by moving Z towards C and Y towards B while maintaining |XZ| = |XY|, since $\alpha = \langle CAB < 90^{\circ}$. If ABC is acute, then

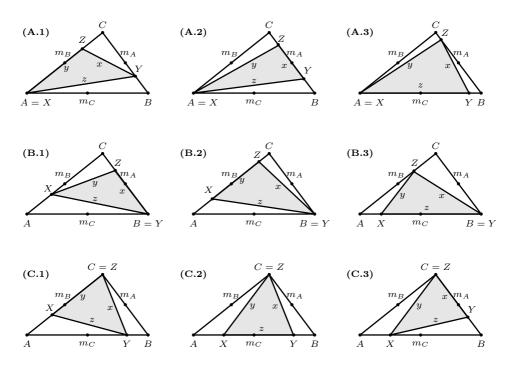


Fig. 7. The 9 possible arrangements of the vertices X, Y, Z in a given triangle ABC.

we can do this until the vertices Z and C will coincide, and the triangle XYZ will be the same as the special embedded triangle $A\overline{B}C$. If ABC is non-acute, then $y \neq z$. Clearly, |AZ| = y > |ZB| > |YZ| = x. Hence, $x \neq y$. A similar argument shows that $x \neq z$.

Subcase A.3: $Y \in AB$ and $Z \in BC$.

Since a < b, the orthogonal projection \widehat{Z} of Z to the line segment AB lies in $m_C B$. If x = y, then $|AY| = 2|A\widehat{Z}| > 2|Am_C| = |AB|$, a contradiction to $Y \in AB$.

If x = z, then the altitude with base z in XYZ is smaller than the altitude with base c in ABC. On the other hand, if $\triangleleft ZYB \ge 90^\circ$, we have x = z < a. In this case, the special embedded triangle A''BC satisfies $\operatorname{area}(A''BC) > \operatorname{area}(XYZ)$. Otherwise, x = z < y (as $\triangleleft AYZ > 90^\circ$) and y < c. Let Y' be the point in AB that is defined by the equality |AY'| = |AZ| (the existence of $Y' \in AB$ is a consequence of y < c). Then, $\operatorname{area}(XY'Z) > \operatorname{area}(XYZ)$. In both cases it follows that the area of XYZ cannot be optimal.

If y = z, consider the special embedded triangle AB'C, define l to be the line parallel to B'C going through Y and let $Z' = l \cap BC$, see Fig. 8. Since $Z' \in CZ$, we have

$$\operatorname{area}(XYZ) < \operatorname{area}(XYZ') = \operatorname{area}(AB'C) \cdot \frac{b + |B'Y|}{b} \cdot \frac{c - b - |B'Y|}{c - b}.$$

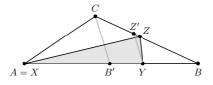


Fig. 8. Illustration for Subcase A.3.

The inequality follows from the fact that $Z' \in CZ$. Therefore, the altitude of XYZ with base z is greater than the altitude of XYZ' with base z. Thus, it is enough to show that

$$\frac{b+|B'Y|}{b}\cdot\frac{c-b-|B'Y|}{c-b}<1.$$

As b > 0 and c - b > 0, this is equivalent to |B'Y|(2b - c + |B'Y|) > 0, which follows from the triangle inequality c < a + b < 2b.

Case B: The common vertex of ABC and XYZ is B = Y.

Subcase B.1: $X \in AC$ and $Z \in BC$.

Since a < c, we have that $\triangleleft AXY > 90^{\circ}$, and hence, we can rotate the triangle XYZ about Y so that the image of the vertices X, Z will be inside of ABC. As in Subcase A.1, this implies that the area of XYZ is not optimal.

Subcase B.2: Both X and Z are in AC.

Observe that b < c implies that A and C are on the same side of the perpendicular bisector of BC. This implies that |XY| = z > |XC| > |XZ| = y. If x = z, we can 'open' $\triangleleft XYZ$ as in Subcase A.2 and get that $\operatorname{area}(XYZ) < \operatorname{area}(\overline{\overline{A}BC})$. Hence, we can assume that x = y.

If the triangle ABC is non-acute, then consider the special embedded triangle ABC_1 . Since the altitudes of ABC_1 and XYZ from vertex B = Y are equal, and $x = y < |BC_1| = |AC_1|$ (as $\triangleleft BCA \ge 90^\circ$), we have that $\operatorname{area}(XYZ) < \operatorname{area}(ABC_1)$.

If ABC is acute, let \hat{B} denote the orthogonal projection of B onto AC. If $Z \in A\hat{B}$, then we can slightly rotate XYZ about Y (as $\triangleleft YXA > \triangleleft YZA > 90^{\circ}$). Thus, by Lemma 13(iv), the area of XYZ is not maximal. Thus, we can assume that $Z \in C\hat{B}$, that is, $\triangleleft YZA \leq 90^{\circ}$. Similarly as above, this implies that x = |YZ| < a = |BC| and thus the special embedded triangle A'BC satisfies $\operatorname{area}(XYZ) < \operatorname{area}(A'BC)$.

Subcase B.3: $X \in AB$ and $Z \in AC$.

If y = z, then, since $\triangleleft CAB < \min(\triangleleft AXZ, \triangleleft ZXY)$, we get that y = |XZ| < |AZ| < b = |AC|, which immediately implies that the special embedded triangle AB'C satisfies area $(XYZ) < \operatorname{area}(AB'C)$.

Now we assume that x = z. If A and Z lie on the same side of the perpendicular bisector of AB, then we can reflect XYZ to this perpendicular bisector.

We denote this reflection by X'Y'Z'. Clearly, $X', Y' \in AB$, and Z' is inside of ABC, which implies that $\operatorname{area}(XYZ)$ is not maximal. If Z is on the perpendicular bisector of AB, then XYZ is strictly contained in the special embedded triangle ABC_1 , so $\operatorname{area}(XYZ) < \operatorname{area}(ABC_1)$. If Z and C are on the same side of the perpendicular bisector of AB, then x = z < |AZ| < |AC| = b, and hence $\operatorname{area}(XYZ) < \operatorname{area}(AB'C)$.

It remains to handle the case x = y. We show that $\operatorname{area}(XYZ) < \operatorname{area}(ABC_1)$. The condition x = y implies that $Z \in C_1C$. Plainly, z = c - |AX|. Denote the lengths of the altitudes from C_1 in ABC_1 and from Z in XYZ by h_{C_1} and h_Z , respectively. Clearly, we get $h_Z = h_{C_1} \frac{c + |AX|}{c}$, and hence

$$\operatorname{area}(XYZ) = \operatorname{area}(ABC_1)\frac{c + |AX|}{c} \cdot \frac{c - |AX|}{c} < \operatorname{area}(ABC_1).$$

Case C: The common vertex of ABC and XYZ is C = Z.

Subcase C.1: $X \in AC$ and $Y \in AB$.

If Y and B are on the same side of the altitude from C, then we can rotate XYZ about Z so that X and Y get to the interior of ABC which by Lemma 13(iv) implies that XYZ is not optimal. If Y and B are on different sides of the altitude from C, then XYZ is strictly contained in the special embedded triangle AB'C.

Subcase C.2: Both X and Y are contained in AB.

If x = y, then we can 'open' $\triangleleft YZX$, which increases its area, thus area(XYZ) is not maximal. Suppose that x = z. If Y and A are on the same side of the altitude from C, then XYZ is strictly contained in the special embedded triangle AB'C. If Y and A lie on different sides of the altitude, then the special embedded triangle A''BC satisfies area $(XYZ) < \operatorname{area}(A''BC)$. Indeed, their altitudes from C are the same, and for their bases we have x = z < a. Thus XYZ is not maximal. A similar argument shows that if y = z, then we have $\operatorname{area}(XYZ) < \operatorname{area}(AB'C)$.

Subcase C.3: $X \in AB$ and $Y \in BC$.

We can rotate XYZ about Z such that the images of X and Y lie in the interior of ABC, and so, by Lemma 13(iv), we get that $\operatorname{area}(XYZ)$ is not maximal.

We have shown that none of the triangles XYZ of the 9 cases in Fig. 7 is a maximum area embedded isosceles triangle of ABC, which completes the proof of Theorem Lemma 3(ii).

4. Maximum Perimeter Embedded Isosceles Triangles — Proof of Lemma 3(iii)

In this section, we prove that for any triangle ABC, any maximum perimeter isosceles triangle XYZ embedded in ABC shares a vertex and the angle at that vertex with ABC. First we collect the observations in Lemmas 5 and 6 concerning maximum perimeter embedded isosceles triangles.

Lemma 15. Let XYZ be a maximum perimeter isosceles triangle embedded in ABC. Then

- (i) each side of ABC contains a vertex of XYZ;
- (ii) no vertex of the triangle XYZ lies in the interior of the triangle ABC;
- (iii) there is a side of ABC which contains a side of XYZ;
- (iv) ABC and XYZ share a vertex.

We will show that an isosceles triangle embedded in ABC which does not share an angle with ABC cannot be of minimum perimeter. Notice that if ABC and XYZ share at least two vertices, then, by Lemma 15(ii), they also share an angle, so we are done. Thus, it is enough to consider those cases where the triangles XYZand ABC share exactly one vertex, without loss of generality the common vertex is A.

Note that in this section, we do not assume a special labeling of ABC, in particular, we do not necessarily have |BC| < |AC| < |AB|. On the other hand, we assume that XYZ is labeled so that |XY| = |YZ|.

We consider the following cases, separately:

Case A: X and Z lie on the same side of ABC.

We can always rotate X or Z (for simplicity, assume it is X) about Y so that the rotated point X' lies in the interior of ABC and $\triangleleft XYZ < \triangleleft X'YZ$, see Fig. 9. By the Hinge theorem (which states that if XYZ and X'Y'Z' are triangles such that XY = X'Y', YZ = Y'Z', and $\triangleleft XYZ < \triangleleft X'Y'Z'$, then $\operatorname{per}(XYZ) < \operatorname{per}(X'Y'Z')$), we get that $\operatorname{per}(XYZ) < \operatorname{per}(X'YZ)$.

Case B: X and Z lie on different sides of ABC.

We will make use of the following classical lemma on the perimeter of the Minkowski sum of convex bodies.

Lemma 16 (see e.g.[29, exercise 4–7]). Let K_1 and K_2 be two convex bodies in the plane and let $K = \frac{K_1+K_2}{2}$ be the Minkowski mean of K_1 and K_2 . Then the perimeter of K is equal to the arithmetic mean of the perimeters of K_1 and K_2 . If K_1 and K_2 are not homothetic triangles, then K is a convex polygon with at least four sides.

The idea is to show that the triangle XYZ is strictly contained in the Minkowski mean M of two other non-homothetic isosceles triangles embedded in ABC, thus,

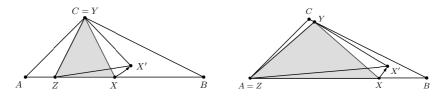


Fig. 9. Illustration for Case A.

by Lemma 16, one of these two must have a strictly larger perimeter (by the fact that if C_1, C_2 are two convex planar sets such that $C_1 \subseteq C_2$, then per $(C_1) \leq \text{per} (C_2)$ [2, 12.10.2]). We proceed by defining 3 subcases and finding such a pair of iscosceles triangles in each subcase separately.

Subcase B.1: The common vertex of ABC and XYZ is A = Y.

If none of X and Z is on the side opposite to Y, then XYZ and ABC have a common angle at Y. Thus, we can assume that either X or Z is on the side opposite to Y, say it is X.

Let δ be a constant satisfying $\delta < \min\{|XB|, |XC|\}$. Define the points X^1 and X^2 by translating X by δ towards C and B, respectively. Let Z^1 and Z^2 be such that they are contained on the side AB with $|YZ^1| = |YX^1|$ and $|YZ^2| = |YX^2|$, see Fig. 10. Let M be the Minkowski mean of X^1YZ^1 and X^2YZ^2 . The vertex Y is contained in both triangles, thus it is also contained in M. It is also easy to see that $X \in M$ since $X = \frac{1}{2}(X^1 + X^2)$. We show that Z is contained in the segment between Y and $\frac{1}{2}(Z^1 + Z^2)$, which implies $Z \in M$. To this end, observe that the segment YX is a median of the triangle X^1YX^2 and thus $|YX| < \frac{1}{2}(|YX^1| + |YX^2|)$, which directly gives that $|YZ| < \frac{1}{2}(|YZ^1| + |YZ^2|)$.

Subcase B.2: The common vertex of ABC and XYZ is A = Z and both X and Y are in the interior of the side of ABC opposite to Z.

Define the points X^1 and X^2 by translating X by δ towards C and B, respectively. We choose δ to be small enough such that there are points Y^1, Y^2 in the segment BCwith $|Y^1Z| = |Y^1X^1|$ and $|Y^2Z| = |Y^2X^2|$, see Fig. 11. Let M be the Minkowski mean of X^1YZ^1 and X^2YZ^2 . As before, it is clear that the vertices X and Z are contained in M.

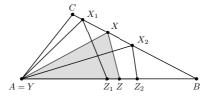


Fig. 10. Illustration for Subcase B.1.

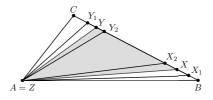


Fig. 11. Illustration for Subcase B.2.

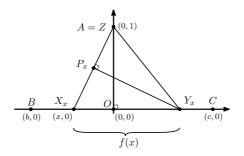


Fig. 12. Embedding the instance of Subcase B.2 in \mathbb{R}^2 .

To argue that $Y \in M$, we shall show that Y is contained in the segment between X and $\frac{1}{2}(Y^1 + Y^2)$. To simplify the calculations, we move and scale the triangle so that A = (0, 1), B = (b, 0), C = (c, 0) and X = (x', 0) with b < x' < c. Note that since $\langle ZXY \rangle$ is acute, x' < 0. For each b < x < 0, let $X_x = (x, 0)$ and Y_x be the point in BC such that $|ZY_x| = |Y_xX_x|$ and define $f(x) = |X_xY_x|$, see Fig. 12. Observe that Y is contained in the segment between X and $\frac{1}{2}(Y^1 + Y^2)$ if and only if $\frac{1}{2}(f(x'-\delta)+f(x'+\delta)) > f(x')$. Thus, it is sufficient to show that f(x) is a convex function on (b, 0).

To find an analytic formula for f(x), we introduce some auxiliary points. Let O = (0,0) and P_x be the orthogonal projection of Y_x to the segment X_xZ . Note that P_x is the midpoint of XZ. Then the triangles $X_xP_xY_x$ and X_xOZ are similar, which yields

$$f(x) = |Y_x X_x| = |X_x Z| \cdot \frac{|X_x P_x|}{|X_x O|} = \sqrt{1 + x^2} \cdot \frac{\sqrt{1 + x^2}/2}{-x} = \frac{1 + x^2}{-2x}$$

The second derivative of f is $f''(x) = -1/x^3$, thus f(x) is convex on the interval (b, 0), which implies that Y is contained in the segment between X and $\frac{1}{2}(Y^1+Y^2)$.

Subcase B.3: The common vertex of ABC and XYZ is A = Z and X, Y lie in the interior of different sides of ABC.

Firstly, since X and Z lie on different sides of ABC, we get that X is on the side opposite to Z, see Fig. 13. If $\triangleleft AXB$ is obtuse, then we can rotate the triangle XYZ about Z and obtain a copy of XYZ which has two vertices in the interior of ABC, thus by Lemma 15, XYZ cannot be of maximum perimeter. Therefore, $\triangleleft AXB$ and consequently $\triangleleft ACB$ are acute.

Define the points X^1 and X^2 by translating X by δ towards C and B, respectively. We choose an increment $\delta \in (0, 1/c)$ which is small enough that there are points Y^1, Y^2 in the segment AC with $|Y^1Z| = |Y^1X^1|$ and $|Y^2Z| = |Y^2X^2|$. Let M be the Minkowski mean of X^1YZ^1 and X^2YZ^2 . The vertices X and Z are clearly contained in M.

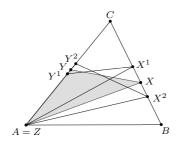


Fig. 13. Illustration for Subcase B.3.

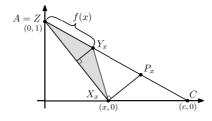


Fig. 14. Embedding the instance of Subcase B.3 in \mathbb{R}^2 .

To prove that $Y \in M$, we shall show that Y is contained in the segment between Z and $\frac{1}{2}(Y^1 + Y^2)$. Again, we translate and scale of the triangle so that A = (0, 1), B = (b, 0), C = (c, 0) and X = (x', 0) with b < x' < c. Since $\triangleleft AXB$ is acute, we have $x' \ge 0$. For each $x \in [x' - \delta, x' + \delta]$, let $X_x = (x, 0)$ and Y_x be the point in ZC such that $|ZY_x| = |Y_xX_x|$ and define $f(x) = |ZY_x|$, see Fig. 14. Note that since $x' \ge 0$ and δ is smaller than 1/c, each $x \in [x' - \delta, x' + \delta]$ satisfies -1/c < x. We want to show that the function f(x) is convex, which then directly implies that Y is contained in the segment between Z and $\frac{1}{2}(Y^1 + Y^2)$.

Let $P_x = (p_1(x), p_2(x))$ be a point on AC such that the segment $P_x X_x$ is orthogonal to AX_x . Note that P_x satisfies $|ZP_x| = 2f(x)$.

Let γ denote the angle $\triangleleft ACX_x$, then $p_2(x) = 1 - 2\sin(\gamma) \cdot f(x)$ which is concave if and only if f(x) is convex. Since $X_x P_x$ is orthogonal to AX_x and P_x is contained in AC, we get the following equations on $p_1(x)$ and $p_2(x)$

$$p_1(x) \cdot x - p_2(x) = x^2$$
, $p_1(x) + cp_2(x) = c$,

which gives $p_2(x) = \frac{cx-x^2}{cx+1}$. Taking the second derivative, we get

$$p_2''(x) = -\frac{2(1+c^2)}{(1+cx)^3} < 0 \text{ for all } x \in \left(-\frac{1}{c}, \infty\right).$$

We showed that none of the triangles of types A and B.1–B.3 can be a maximum perimeter embedded isosceles triangle of ABC, which completes the proof of Theorem 3(iii).

5. Minimum Perimeter Enclosing Triangles — Proof of Theorem 4

In this section, we prove that any smallest perimeter isosceles container of a triangle is either a special container or one of two non-special containers defined in the next subsection. We also show that this is the shortest possible characterization of isosceles containers, that is, any of the five examples is realized as a minimum perimeter isosceles container for some triangle ABC. Now, we define two non-special isosceles containers that can be optimal.

Note that in this section, we do not assume a special labeling of ABC, in particular, we do not necessarily have |BC| < |AC| < |AB|. Furthermore, we assume that the isosceles containers of ABC are labelled with PRS satisfying |PR| = |RS|.

5.1. Two examples for non-special minimum perimeter containers of a triangle

Let P be a point in \mathbb{R}^2 and l a line such that $P \notin l$ and let m denote the distance of P from l. Define an isosceles triangle $PR^{\gamma}S^{\gamma}$ such that S^{γ} and R^{γ} lie on l and its apex angle γ is in R^{γ} , see Fig. 15.

Proposition 17. The perimeter function $p(\gamma) = per(PR^{\gamma}S^{\gamma})$ has a unique minimum at

$$\gamma^* = 4 \tan^{-1} \left(\frac{1}{2} (1 + \sqrt{5} - \sqrt{2(1 + \sqrt{5})}) \right) \approx 76.3466^{\circ}.$$
 (1)

Proof outline. It is easy to see that $|PR^{\gamma}| = |R^{\gamma}S^{\gamma}| = \frac{m}{\sin \gamma}$ and $|PS^{\gamma}| = \frac{m}{\sin(90^{\circ} - \gamma/2)} = \frac{m}{\cos(\gamma/2)}$. Hence per $(PR^{\gamma}S^{\gamma}) = m\left(\frac{2}{\sin \gamma} + \frac{1}{\cos(\gamma/2)}\right)$. Elementary analysis shows that the function $f(x) = \frac{2}{\sin x} + \frac{1}{\cos(x/2)}$ is strictly decreasing in $(0^{\circ}, \gamma^{*}]$ and strictly increasing in $[\gamma^{*}, 180^{\circ})$. Thus it has a unique minimum in $0 \le x \le 180^{\circ}$ that is taken at the value specified in Eq. (1).

Example 18. Let $PRS = PR^{\gamma^*}S^{\gamma^*}$ be an isosceles triangle with apex angle γ^* which is defined as in Proposition 17. Let ABC be an acute triangle in PRS such that ABC and PRS have exactly one common vertex at A = P and $B, C \in SR$ (see Fig. 16). Furthermore, ABC is such that the largest angle γ of ABC is at C with $\gamma < \gamma^*$ being close to γ^* (e.g., 76°) and ABC is almost isosceles ($|AC| \approx |BC|$).

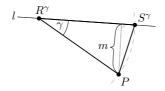


Fig. 15. Illustration for Proposition 17.

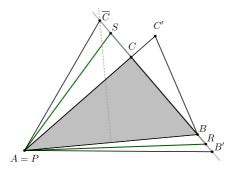


Fig. 16. Illustration for Example 18.

Claim 1. The perimeter of PRS is strictly smaller than the perimeter of any special container of ABC.

Proof outline. By Corollary 11, it is enough to show that the special containers AB'C, ABC', and $AB\overline{C}$ have larger perimeter than PRS. First observe that, since $a \approx b < c$, ABC is an 'almost' isosceles triangle, thus the perimeter per $(AB'C) \approx per(ABC)$ and per $(ABC') > d \cdot per(ABC)$, for a fixed d > 1. This implies that per (AB'C) < per(ABC'). Now we show that PRS has perimeter smaller than per (AB'C) and per $(AB\overline{C})$. Note that, each of PRS, AB'C and $AB\overline{C}$ are isosceles triangles with base vertex A = P and legs on the line RS. By Proposition 17, the smallest perimeter isosceles triangle under these conditions is PRS. Thus, it is enough to guarantee that the triangles AB'C and $AB\overline{C}$ do not coincide with PRS which follows from the fact that ABC and PRS has exactly one common vertex.

Now we turn to our second example. We start by taking the points A = P = (0,0), C = (1,v) and $S_x = (x,0)$ and define R_x to be the point on the S_xC ray so that $|PR_x| = |R_xS_x|$. The next claim follows by elementary calculations, its proof is omitted.

Proposition 19. For any $x \in (1,2)$, the perimeter of PR_xS_x can be expressed as

per
$$(PR_xS_x) = f_v(x) = x\left(1 + \sqrt{1 + \frac{v^2}{(1-x)^2}}\right).$$
 (2)

and for any $v \in [0.56, \sqrt{3})$, the function f_v has a unique minimum in (1, 2) denoted by x_v^* .^a

Example 20. Consider a triangle ABC that can be embedded in \mathbb{R}^2 as A = (0,0), C = (1,v) and $B = (x_b,0)$ with $1 < x_b < x_v^*$ (the value x_v^* is defined in

^aThe formula for x_v^* is given in [1].

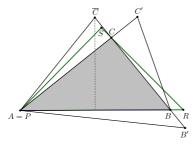


Fig. 17. Illustration for Example 20.

Proposition 19; see also Fig. 17). Let PRS be the an isosceles triangle with P = A, $S = (x_v^*, 0)$, and R defined as the point on the S_xC ray with |PR| = |RS|. By definition, SPR is an isosceles container of ABC.

Claim 2. The perimeter of PRS is smaller than the perimeter of any special container of ABC.

Proof outline. By Corollary 11, we only need to show that *PRS* has a smaller perimeter than the special containers AB'C, ABC', and $AB\overline{C}$. Observe that by the choices of x_v^* and x_b , we have per $(AB\overline{C}) = f_v(x_b) < f_v(x_v^*) = \text{per}(PRS)$.

We verify the remaining cases only for the fixed value v = 0.7. The function $f_{0.7}(x)$ takes its minimum at $x^*_{0.7} \approx 1.57517$, and thus per $(PRS) = f_{0.7}(x^*_{0.7}) \approx 4.056333$. On the other hand, if we set e.g. $x_b = 1.57$, we have per $(AB'C) \approx 4.229145$ and per $(ABC') \approx 4.084007$.

5.2. Proof of Theorem 4

We start by proving that every smallest perimeter isosceles container of a triangle $\Delta = ABC$ is either a special container or one of the two triangles constructed in the Examples 18 and 20. Later, we will show that each of these five containers is realized as the unique minimum perimeter isosceles container for some triangle ABC. By Lemmas 5 and 6, we have the following statements on minimum perimeter isosceles containers.

Lemma 21. Let PRS be any minimum area isosceles triangle enclosing ABC. Then

- (i) a side of PRS contains a side of ABC;
- (ii) each side of PRS contains a vertex of ABC;
- (iii) ABC and PRS share a common vertex;
- (iv) no vertex of ABC lies in the interior of PRS.

If PRS shares the vertex R with ABC, but it does not share the angle at R, then we can get a smaller perimeter container by decreasing $\triangleleft SRP$ (while keeping

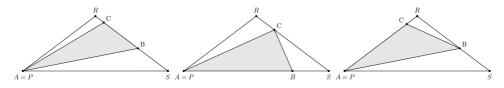


Fig. 18. Illustration for Case 1 (left and middle) and Case 2 (right).

|PR| = |RS| unchanged). Thus without loss of generality, we can assume that PRS shares the vertex P with ABC. The above restrictions allow only the following two types of minimum perimeter isosceles containers that do not share an angle with ABC (see also Fig. 18):

Case 1: If two vertices of ABC lie in the interior of RS, or one of the vertices of ABC lies in the interior of the side RS and one lies in the interior of PS.

The smallest perimeter isosceles containers of these types are precisely the nonspecial optimal containers shown in Examples 18 and 20.

Case 2: One vertex of ABC is in the interior of PR and one is in the interior of RS.

Let T denote the base of the altitude perpendicular to RS and let B denote the vertex in RS. If $|SB| \leq |ST|$, then $\triangleleft SBP \geq 90^{\circ}$, hence we can rotate ABC about A = P such that the triangle remains in PRS and hence PRS was not minimal, see Fig. 19. Note that this happens if PRS is not acute. From now on, we assume that $\triangleleft SBP < 90^{\circ}$, which implies |AB| < |AR| if $B \neq R$.

If |AC| < |AB|, then we take $C' \in AR$ such that |AB| = |AC'| < |AR| so $AC' \subset AR$ (Fig. 19). Thus, ABC' is an isosceles container of ABC and $ABC' \subsetneq PRS$. Hence, PRS was not minimal. Therefore, we may assume that |AC| > |AB|, as |AC| = |AB| would imply that ABC was isosceles.

If $\triangleleft RAB < \triangleleft BRA$ holds, let B' be the point on the line AB such that |AC| = |AB'| then we have |AB'| = |AC| < |AR| = |RS|, and hence per (AB'C) <

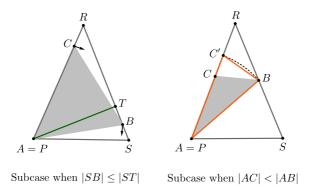


Fig. 19. Simple configurations of Case 2.

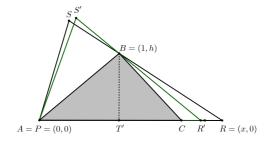


Fig. 20. Case 2 embedded in \mathbb{R}^2 .

per (PRS), thus PRS was not minimal. Thus, assume that $\triangleleft BRA \leq \triangleleft RAB$ and as $\triangleleft RAB + \triangleleft BRA = \triangleleft SBP < 90^{\circ}$, we get that $\triangleleft BRA = \triangleleft PRS < 45^{\circ}$.

For the remaining part, we embed the configuration in \mathbb{R}^2 such that P = A = (0,0), R = (x,0) and B = (1,h), where x > 1 and h > 0, see Fig. 20. Under the assumptions that |AC| > |AB| and $\triangleleft PRS < \min(\triangleleft RAB, 45^\circ)$, we show that per (PRS) as a function of x is increasing. Thus, as $B \neq R, C \neq R$ there exists a smaller perimeter isosceles container of ABC than PRS (e.g., PR'S' in Fig. 20). The condition $\triangleleft PRS < \triangleleft RAB$ implies that |PT'| < |RT'|, where T' is the base of the altitude of PR, hence x > 2.

Clearly,
$$|BR| = \sqrt{h^2 + (x-1)^2}$$
 and $\sin(\triangleleft PRS) = \frac{h}{\sqrt{h^2 + (x-1)^2}}$. Hence
per $(PRS) = 2x(1 + \sin(\frac{\triangleleft PRS}{2}))$. As $\sin \delta = 2\sin(\frac{\delta}{2})\sqrt{1 - \sin^2(\frac{\delta}{2})}$, we get

$$\sin\left(\frac{\triangleleft PRS}{2}\right) = \frac{1}{\sqrt{2}}\sqrt{1 \pm \frac{x-1}{h}\sqrt{\frac{1}{1+\left(\frac{x-1}{h}\right)^2}}},$$

where the \pm is taken to be a - sign, since $\triangleleft PRS < 45^{\circ}$. Therefore,

per
$$(PRS) = 2x + \sqrt{2}x\sqrt{1 - \frac{x-1}{h}\sqrt{\frac{1}{1 + (\frac{x-1}{h})^2}}}.$$

Let $y = \frac{x-1}{h}$ and let $f_h(y) = (1 + hy)(1 + \sqrt{\frac{1-y}{2}\sqrt{\frac{1}{1+y^2}}})$. It follows from our assumptions that y > 1/h. We show that $f_h(y)$ is strictly increasing in y, which implies that PRS is not a minimum perimeter isosceles container of ABC. For $g(y) := 1 + \sqrt{\frac{1-y}{2}\sqrt{\frac{1}{1+y^2}}}$, we show that $f'_h(y) = ((1 + hy)g(y))' > 0$, equivalently $-g'(y) < \frac{hg(y)}{1+hy}$. Simple calculation shows that

$$-g'(y) = \frac{1}{2\sqrt{2}(1+y^2)} \sqrt{1 + \sqrt{\frac{y^2}{1+y^2}}} < \frac{1}{2(1+y^2)},$$

where the last inequality holds as $\frac{y^2}{1+y^2} < 1$ for all $y \in \mathbb{R}$. Note that g(y) > 1, hence hg(y) > h. Thus, it is enough to show that

$$\frac{1}{2(1+y^2)} < \frac{h}{1+hy} \quad \text{if } y = \frac{x-1}{h} > \frac{1}{h}$$

This is true if and only if $0 < 2hy^2 - hy + 2h - 1$, which holds if its roots satisfy $y_1 < y_2 = \frac{h + \sqrt{-15h^2 + 8h}}{2h} \le \frac{1}{h}$. The last inequality is equivalent to $0 \le 4h^2 - 3h + 1 = 1$ $(2h-1)^2 + h$, which is true for h > 0. Therefore, the argument above verifies that in this case PRS is not minimal. This concludes the proof in Case 2.

Note on realizability. Now we briefly discuss that each of the special containers AB'C, ABC', $AB\overline{C}$, and triangles constructed in Examples 18 and 20 can occur as a minimum perimeter container for some ABC. It is easy to find triangles for which one of the special containers is the best among the five options.

To see that the container of Example 20 is optimal for some triangles, note that the construction presented in Example 18 works only if the special containers of ABC satisfy $\gamma^* \in (\triangleleft AB'C, \triangleleft AB\overline{C})$. Now consider the example from the proof of Claim 2. It can be easily calculated that under these choices $\triangleleft(BCA) \approx 78,310868^{\circ}$.

This (together with Claim 2) implies that for the example presented in the proof of Claim 2, the container described in Example 20 is better than the one given in Example 18 and than any special container.

Finally, we show that the container presented in Example 18 is optimal for some triangles. Following the construction in the proof of Claim 2, consider the triangle ABC with A = (0,0), and C = (1,0.8) and $B = (0, x_{0.8}^*)$ such that $f_{0.8}(x)$ takes its minimum at $x_{0.8}^* \approx 1.62474$. We get that the container constructed in Example 20 coincides with the special container ABC and per (ABC) = $f_{0.8}(x_{0.8}^*) \approx 4.264511$. Simple calculation shows that per $(ABC') \approx 4.3250804$, thus per $(AB\overline{C}) < \text{per}(ABC')$. Since $\triangleleft(B\overline{C}A) \approx 75.974334^\circ < \gamma^* < \triangleleft(BCA) =$ $\triangleleft(B'CA) \approx 89.327359^{\circ}$, the construction of Example 18 provides smaller perimeter than any of the special containers, indeed if we let SPR to be the container constructed in Example 18 for out choice of ABC, then we get per $(PRS) \approx 4.264431$.

This concludes the proof of Theorem 4.

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