Information Theory Exercises to calculate channel capacity

1. Erasure channel with four input letters.

Let the input alphabet be $\mathcal{X} = \{0, 1, 2, 3\}$ and the output alphabet be $\mathcal{Y} = \{0, 1, 2, 3, *\}$. The transition probabilities are given by W(i|i) = 1 - p and W(*|i) = p for every $i \in \mathcal{X}$, while W(j|i) = 0 for $j \notin \{i, *\}$. Give the capacity of this channel.

Solution. We claim that C = 2(1 - p).

First we show $C \ge 2(1-p)$. This can be done by assuming the input distribution to be uniform and calculating I(X,Y) as follows.

$$I(X,Y) = H(Y) - H(Y|X) =$$

$$H((1-p)/4, (1-p)/4, (1-p)/4, (1-p)/4, p) - h(p) =$$

$$h(p) + (1-p)\log 4 - h(p) = 2(1-p).$$

To prove this is also an upper bound, we introduce the indicator variable for erasure: Let E = 1 if Y = * and E = 0 otherwise. Then we can write

$$H(Y) = H(Y, E) = H(E) + H(Y|E) =$$

$$h(p) + (1-p)H(Y|E = 0) \le h(p) + (1-p)\log 4,$$

where the first equality is by E being determined by Y and the last inequality is implied by the fact that Y can take four values that are different of *, that is, four values when E = 0. We also used that H(Y|E = 1) = 0, since if E = 1 then Y is determined to be *. Thus

$$I(X,Y) = H(Y) - H(Y|X) \le h(p) + (1-p)\log 4 - h(p) = 2(1-p).$$

So 2(1-p) bounds C both from above and from below, thus it is the exact value of the capacity C.

2. Let the input alphabet be $\mathcal{X} = \{0, 1, 2, 3, 4, 5, 6, 7\}$ and let the output alphabet \mathcal{Y} be the same. The transition probabilities are given by W(i|i) = 1-p-q, W(i+1|i) = p, and W(i+2|i) = q for every $i \in \mathcal{X}$, where addition is performed modulo 8. All other transition probabilities are 0. Give the capacity of this channel.

Solution:

$$I(X,Y) = H(Y) - H(Y|X) = H(Y) - H(p,q,1-p-q) \le \log 8 - H(p,q,1-p-q)$$

The right had side can be attained, since by the symmetry of the system uniform input distribution results in uniform output distribution, and then $H(Y) = \log 8 = 3$ holds. Thus C = 3 - H(p.q, 1 - p - q).